

CS41B RECURSION

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CS 52 – Fall 2015

Admin

Assignment 4 out: due Monday (2/29 at 11:59pm)

Survey in assignment 4

Assignment 2 scores

Midterm back

- Average: 23.5 (81%)
- Q1: 26.9 (93%)
- Median: 23.75 (82%)
- Q3: 20.25 (70%)

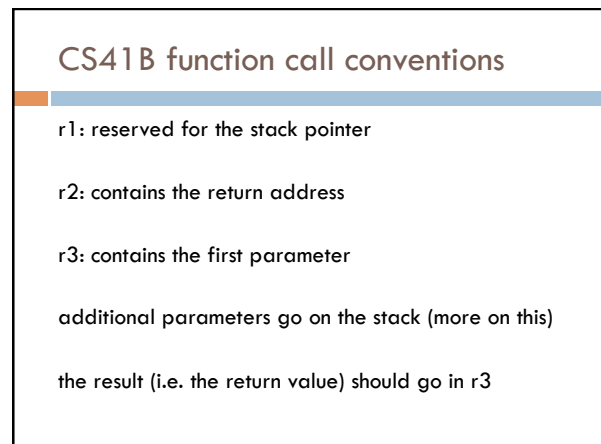
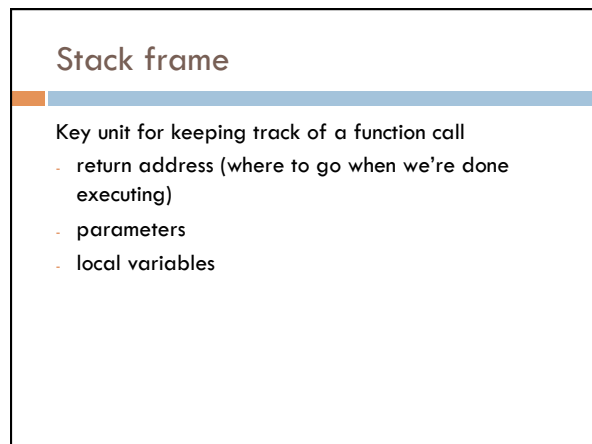
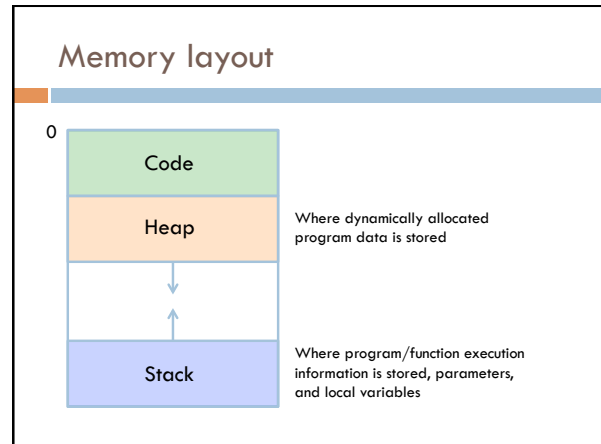
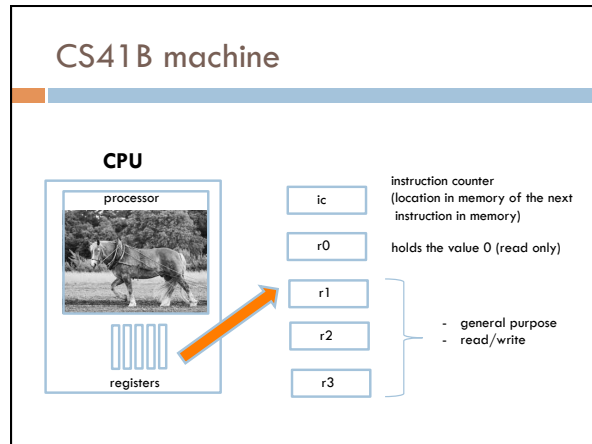
Academic Honesty

A few rules to follow for this course to keep you out of trouble:

- If you talk with someone in the class about a problem, you should not take notes. If you understand the material you talked about, you should be able to recreate it on your own.
- Similarly, if you talk with someone, you must wait 5 minutes before resuming work on the problem. Stretch. Use the restroom. Go for a quick walk. This will ensure that you really understand the material.
- You may not sit next to (or where you can see the screen of) anyone you are talking with about the assignment.
- The only time you may look at someone else's screen is if they are asking you for help with a basic programming problem (e.g. syntax error). You should not look at someone else's code to help yourself!

Examples from this lecture

<http://www.cs.pomona.edu/~dkauchak/classes/cs52/examples/cs41b/>



Structure of a single parameter function

```
fname
    psh r2          ; save return address on stack
    ...            ; do work using r3 as argument
                  ; put result in r3

    pop r2         ; restore return address from stack
    jmp r2         ; return to caller
```

conventions:

- argument is in r3
- r1 is off-limits since it's used for the stack pointer
- return value goes in r3

Our first function call

```
    loa r3 r0      ; get variable

    lclw r2 increment ; call increment
    cal r2 r2

    sto r0 r3      ; write result,
    hlt           ; and halt

increment
    psh r2         ; save the return address on the stack
    adc r3 r3 1    ; add 1 to the input parameter
    pop r2         ; get the return address from stack
    jmp r2         ; go back to where we were called from
```

Functions with multiple arguments

```
fname
    psh r2          ; save return address on stack
    loa r2 r1 4     ; load the second parameter into r2
    ...            ; do work using r3 and r2 as arguments
                  ; put result in r3

    pop r2         ; restore return address from stack
    jmp r2         ; return to caller
```

conventions:

- first argument is in r3
- r1 is off-limits since it's used for the stack pointer
- return value goes in r3

Functions with multiple arguments

```
fname
    psh r2          ; save return address on stack
    loa r2 r1 4     ; load the second parameter into r2
    ...            ; do work using r3 and r2 as arguments
                  ; put result in r3

    pop r2         ; restore return address from stack
    jmp r2         ; return to caller
```

| | | |
|-----|-------|------------------------|
| loa | RR[S] | dest = mem[src0 + arg] |
|-----|-------|------------------------|

What does this operation do? What is the 4?

Functions with multiple arguments

```
fname
    psh r2          ; save return address on stack
    loa r2 r1 4    ; load the second parameter into r2
    ...           ; do work using r3 and r2 as arguments
    ...           ; put result in r3
    pop r2         ; restore return address from stack
    jmp r2         ; return to caller
```

| | | |
|-----|-------|------------------------|
| loa | RR[S] | dest = mem[src0 + arg] |
|-----|-------|------------------------|

- r1 is the stack pointer and points at the top (next) slot
- stacks grow towards smaller memory values

Functions with multiple arguments

```
fname
    psh r2          ; save return address on stack
    loa r2 r1 4    ; load the second parameter into r2
    ...           ; do work using r3 and r2 as arguments
    ...           ; put result in r3
    pop r2         ; restore return address from stack
    jmp r2         ; return to caller
```

| | | |
|-----|-------|------------------------|
| loa | RR[S] | dest = mem[src0 + arg] |
|-----|-------|------------------------|

- r1 is the stack pointer and points at the top (next) slot
- stacks grow towards smaller memory values
- r1+2 is then the top value of the stack
- r1+4 is the 2nd value of the stack

Multiple arguments

```
max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    adc r3 r2 0
endif
    pop r2
    jmp r2
```

What does this code do?

Multiple arguments

```
max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    adc r3 r2 0
endif
    pop r2
    jmp r2
```

max, as a function!

Calling max

```

loa r3 r0
loa r2 r0

psh r2
lcw r2 max
cal r2 r2
pop r2

sto r0 r3
hlt

```

Anything different?

Calling max

```

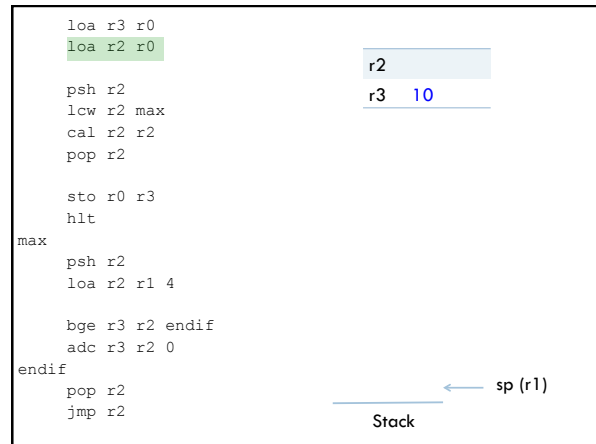
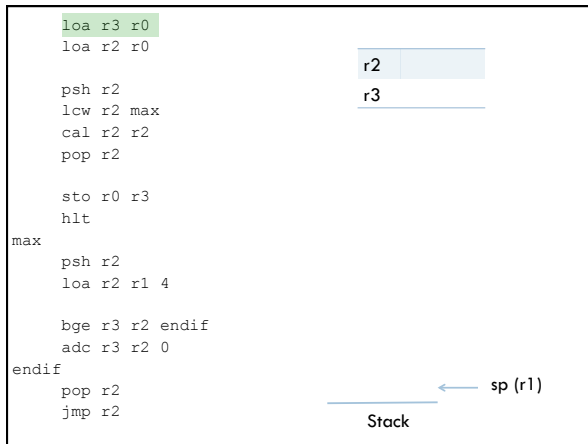
loa r3 r0
loa r2 r0

psh r2
lcw r2 max
cal r2 r2
pop r2

sto r0 r3
hlt

```

For the second argument,
psh it on the stack



```

loa r3 r0
loa r2 r0
psh r2
lcw r2 max
cal r2 r2
pop r2

sto r0 r3
hlt
max
psh r2
loa r2 r1 4

bge r3 r2 endif
adc r3 r2 0
endif
pop r2
jmp r2
    
```

| | |
|----|----|
| r2 | 2 |
| r3 | 10 |

← sp (r1)
Stack

```

loa r3 r0
loa r2 r0
psh r2
lcw r2 max
cal r2 r2
pop r2

sto r0 r3
hlt
max
psh r2
loa r2 r1 4

bge r3 r2 endif
adc r3 r2 0
endif
pop r2
jmp r2
    
```

| | |
|----|----|
| r2 | 2 |
| r3 | 10 |

← sp (r1)
2
Stack

```

loa r3 r0
loa r2 r0
psh r2
lcw r2 max
cal r2 r2
pop r2

sto r0 r3
hlt
max
psh r2
loa r2 r1 4

bge r3 r2 endif
adc r3 r2 0
endif
pop r2
jmp r2
    
```

| | |
|----|----|
| r2 | 2 |
| r3 | 10 |

← sp (r1)
2
Stack

```

loa r3 r0
loa r2 r0
psh r2
lcw r2 max
cal r2 r2
pop r2

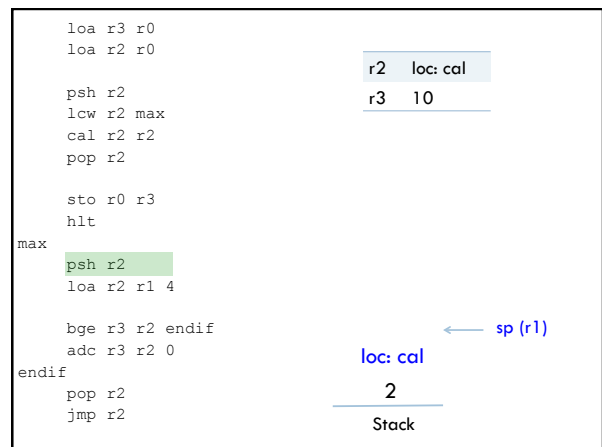
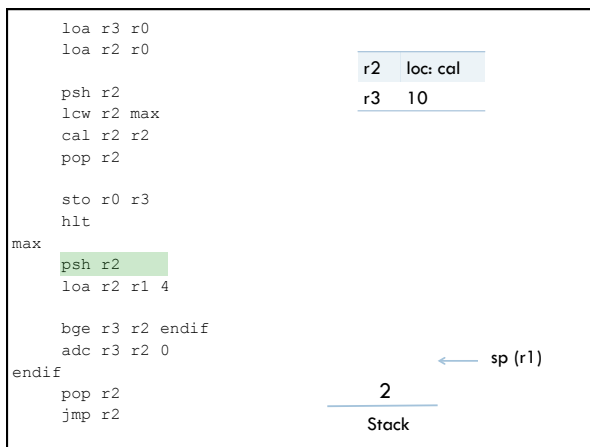
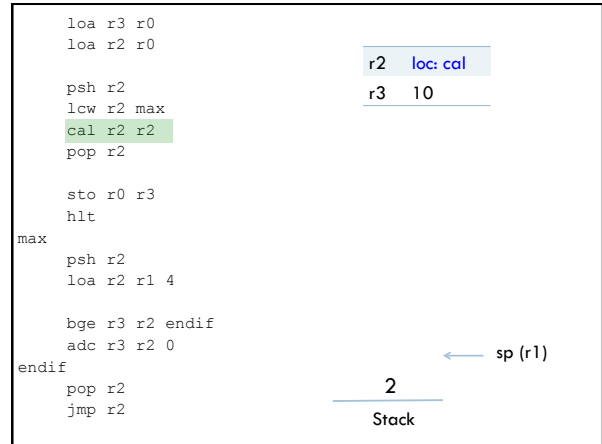
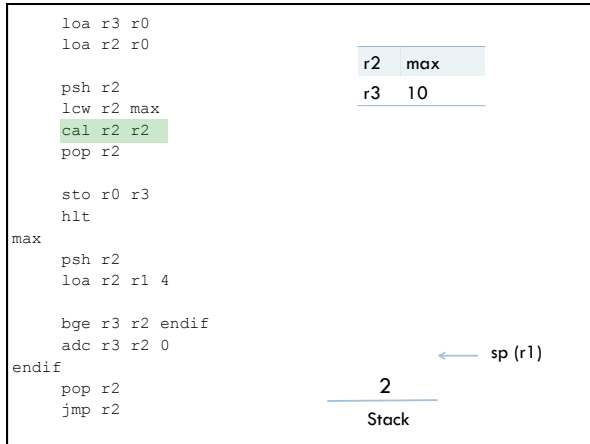
sto r0 r3
hlt
max
psh r2
loa r2 r1 4

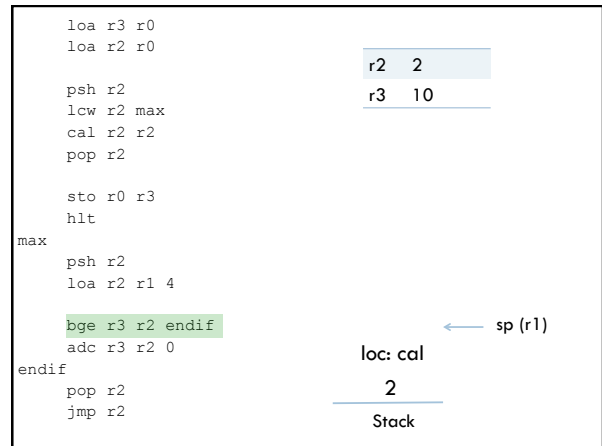
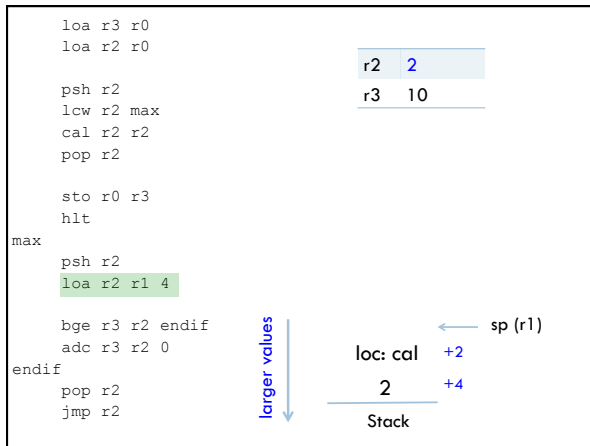
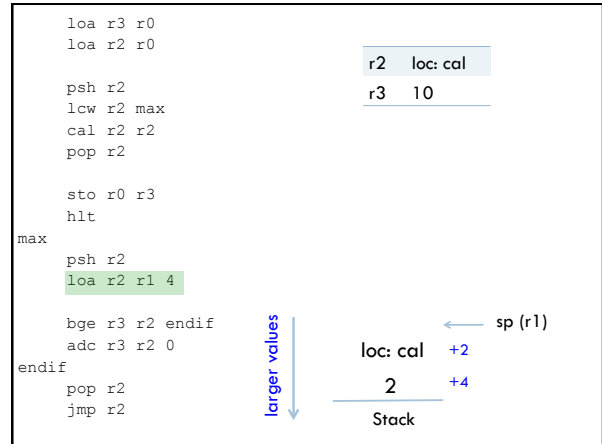
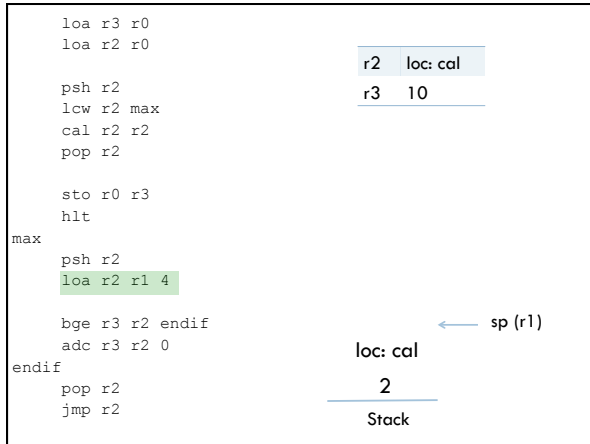
bge r3 r2 endif
adc r3 r2 0
endif
pop r2
jmp r2
    
```

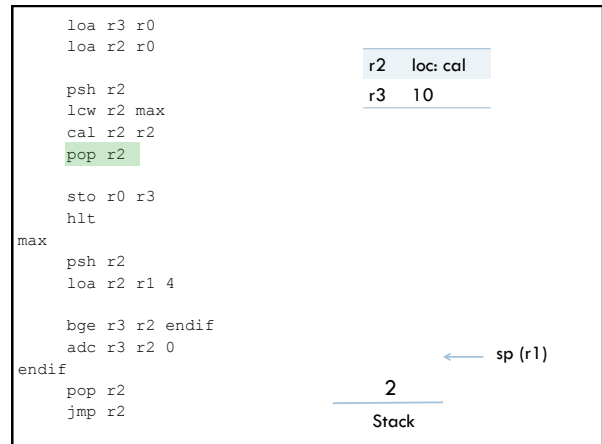
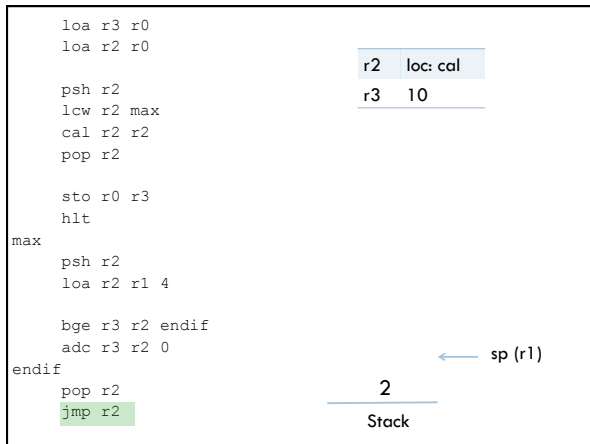
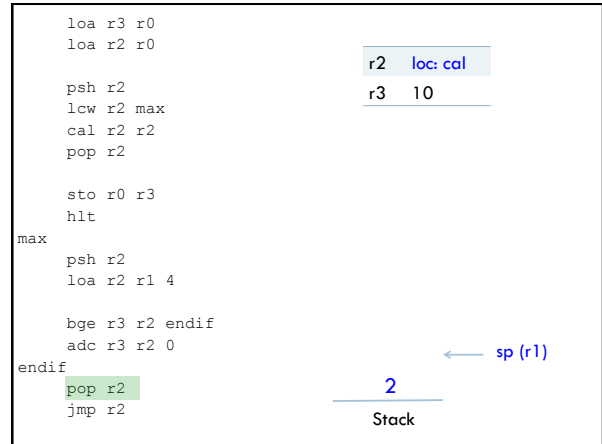
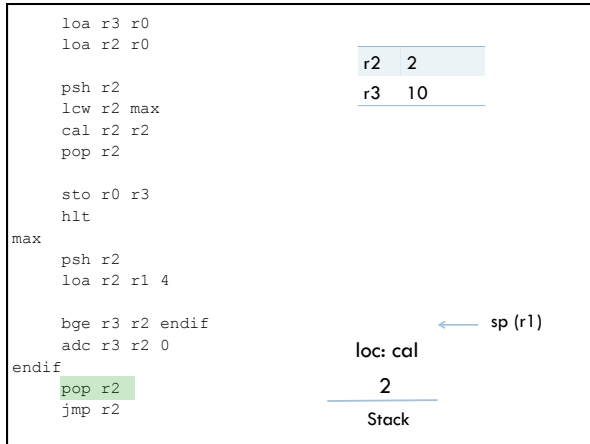
| | |
|----|-----|
| r2 | max |
| r3 | 10 |

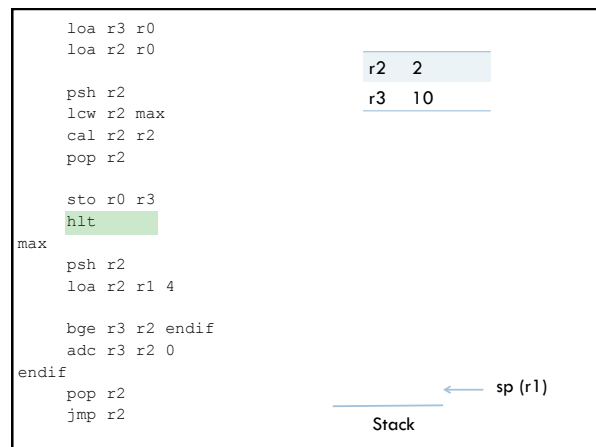
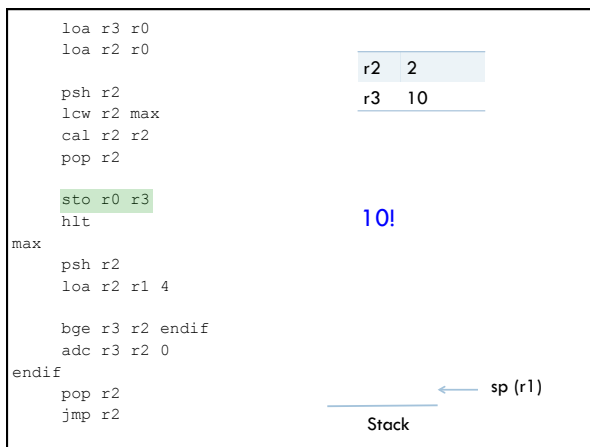
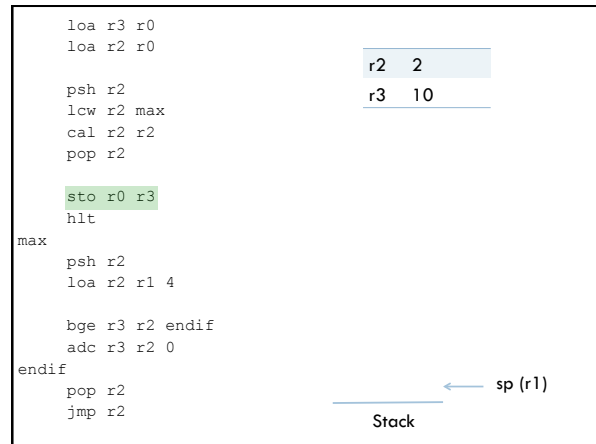
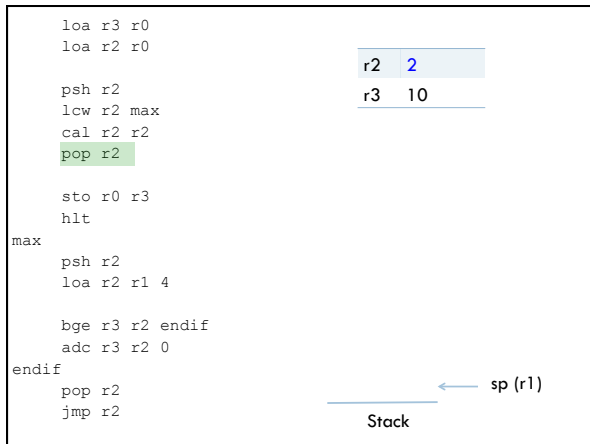
← sp (r1)
2
Stack

Notice that we overwrote the value in r2
If we hadn't saved it on the stack, it would have been lost









Real structure of CS41B program

```

; great comments at the top!
;
    lw r1 stack                Save address of highest end
                                (highest address) of the stack in r1

    instruction1                ; comment
    instruction2                ; comment
    ...
    hlt

;
; stack area: 50 words
;
    dat 100
stack
end

```

Reserve 50 words for the stack

Recursion

```

int mystery(int a, int b){
    if( b <= 0 ){
        return 0
    }
    else
        return a + mystery(a, b-1)
}

```

What does this function do?

Recursion

```

int mystery(int a, int b){
    if( b <= 0 ){
        return 0
    }
    else
        return a + mystery(a, b-1)
}

```

Multiplication... $a*b$ (assuming b is positive)

Note to future Dave from past Dave: write the function up on the board ☺

```

mult
    psh r2                ; save the return address
    loa r2 r1 4           ; get at the 2nd argument, b
                        ; a = r3, b = r2
                        } Function startup

    bgt r2 r0 else       ; r2 > 0, i.e. recursive case
    adc r3 r0 0          ; return 0
    brs endif           } Base case

else
    sbc r2 r2 1          ; r2 = b-1

    psh r3                ; save first argument, a, on stack
                        ; (it's going to get overwritten by the return!)
    psh r2                ; add r2 as 2nd argument, r3 shouldn't have changed
    lw r2 mult           ; call mult recursively
    cal r2 r2             Recursive call
    pop r0                ; pop 2nd argument off stack
                        } Recursive case

    loa r2 r1 2          ; load a into r2 off of the stack
    add r3 r3 r2         ; r3 = a + mult(a, b-1)
    pop r0                ; remove first argument (a) from stack
                        } answer calculation

endif
    pop r2                ; get the return address
    jmp r2                ; return
                        } Function cleanup and return

```

```

mult
  psh r2      ; save the return address
  loa r2 r1 4 ; get at the 2nd argument, b
              ; a = r3, b = r2

  bgt r2 r0 else ; r2 > 0, i.e. recursive case
  adc r3 r0 0   ; return 0
  brs endif

else
  sbc r2 r2 1  ; r2 = b-1

  psh r3      ; save first argument, a, on stack
              ; (it's going to get overwritten by the return!)
  psh r2      ; add r2 as 2nd argument, r3 shouldn't have changed
  lcw r2 mult ; call mult recursively
  cal r2 r2
  pop r0      ; pop 2nd argument off stack

  loa r2 r1 2 ; load a into r2 off of the stack
  add r3 r3 r2 ; r3 = a + mult(a, b-1)
  pop r0      ; remove first argument (a) from stack
endif
  pop r2      ; get the return address
  jmp r2      ; return
  
```

Function startup

if(b <= 0)
return 0

mystery(a, b-1)

a + mystery(a, b-1)

Function cleanup and return

```

mult
  psh r2      ; save the return address
  loa r2 r1 4 ; get at the 2nd argument, b
              ; a = r3, b = r2

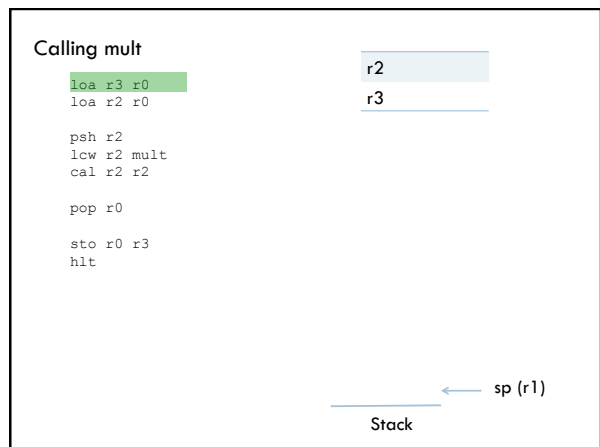
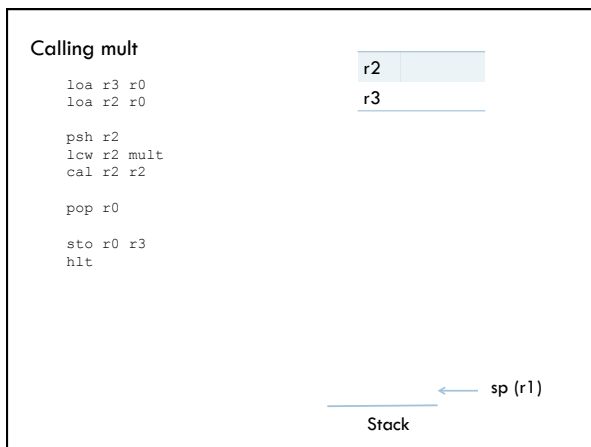
  bgt r2 r0 else ; r2 > 0, i.e. recursive case
  adc r3 r0 0   ; return 0
  brs endif

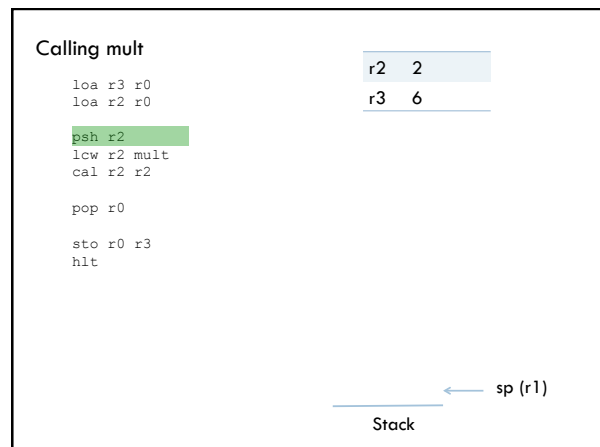
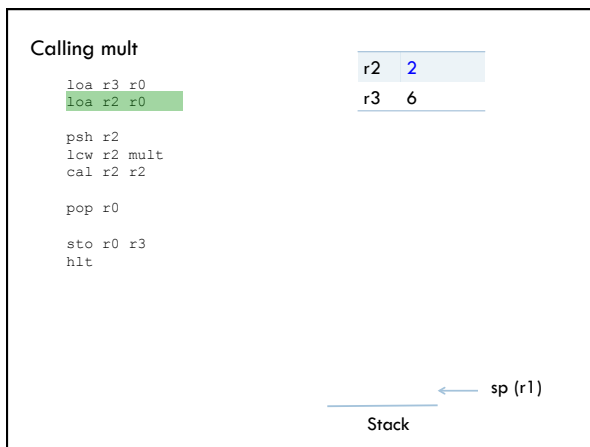
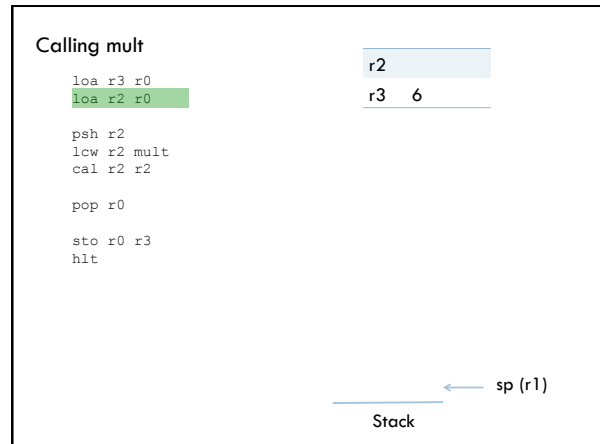
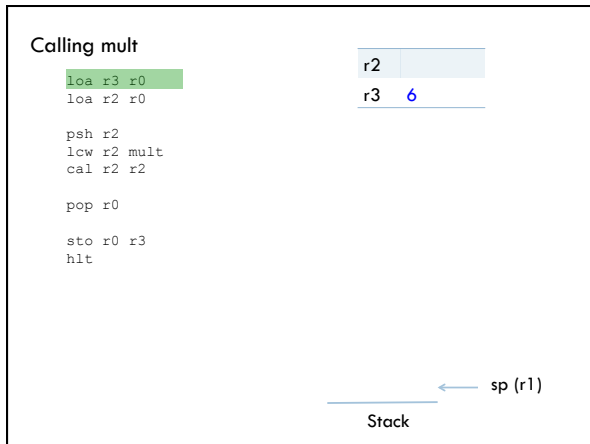
else
  sbc r2 r2 1  ; r2 = a-1

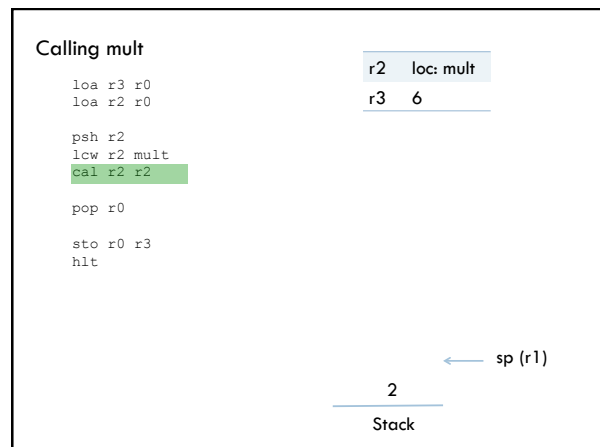
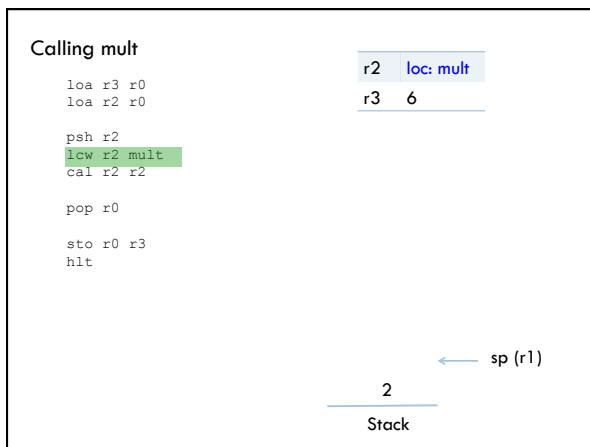
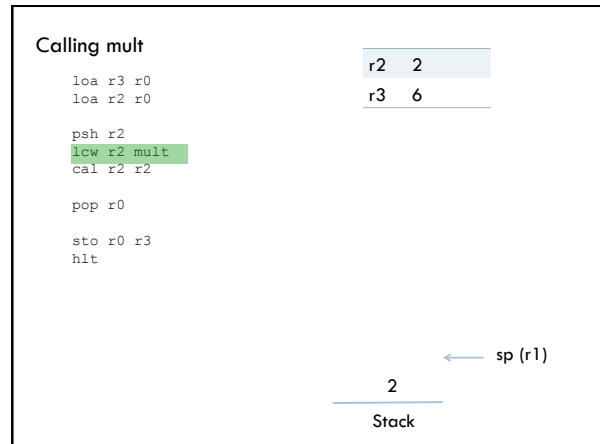
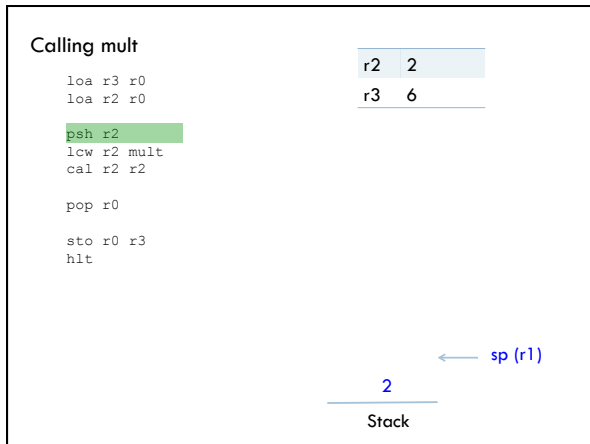
  psh r3      ; save first argument, a, on stack
              ; (it's going to get overwritten by the return!)
  psh r2      ; add r2 as 2nd argument, r3 shouldn't have changed
  lcw r2 mult ; call mult recursively
  cal r2 r2
  pop r0      ; pop 2nd argument off stack

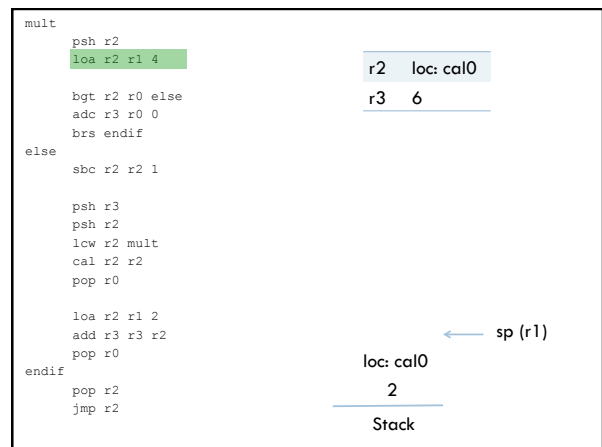
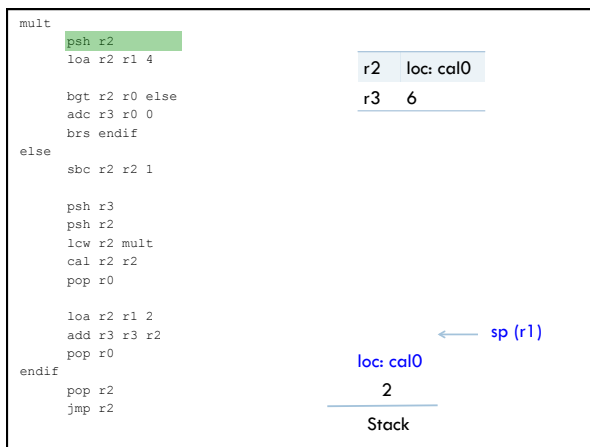
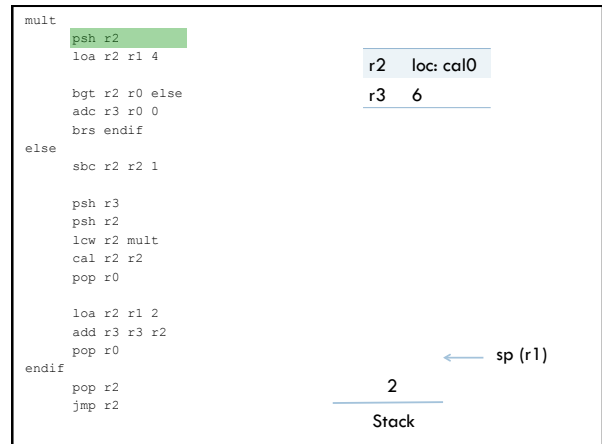
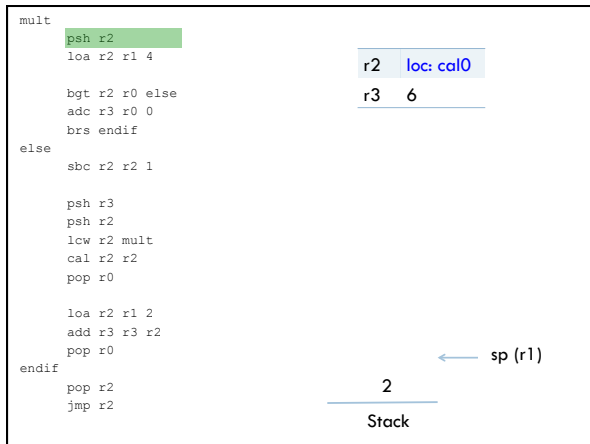
  loa r2 r1 2 ; load a into r2 off of the stack
  add r3 r3 r2 ; r3 = a + mult(a, b-1)
  pop r0      ; remove first argument (a) from stack
endif
  pop r2      ; get the return address
  jmp r2      ; return
  
```

Notice symmetry of psh and pop









```

mult
  psh r2
  loa r2 r1 4
  bgt r2 r0 else
  adc r3 r0 0
  brs endif
else
  sbc r2 r2 1

  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  loa r2 r1 2
  add r3 r3 r2
  pop r0
endif
pop r2
jmp r2
  
```

| | |
|----|-----------|
| r2 | loc: cal0 |
| r3 | 6 |

← sp (r1)

| | |
|-----------|----|
| loc: cal0 | +2 |
| 2 | +4 |

Stack

larger values ↓

```

mult
  psh r2
  loa r2 r1 4
  bgt r2 r0 else
  adc r3 r0 0
  brs endif
else
  sbc r2 r2 1

  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  loa r2 r1 2
  add r3 r3 r2
  pop r0
endif
pop r2
jmp r2
  
```

| | |
|----|---|
| r2 | 2 |
| r3 | 6 |

← sp (r1)

| | |
|-----------|----|
| loc: cal0 | +2 |
| 2 | +4 |

Stack

larger values ↓

```

mult
  psh r2
  loa r2 r1 4
  bgt r2 r0 else
  adc r3 r0 0
  brs endif
else
  sbc r2 r2 1

  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  loa r2 r1 2
  add r3 r3 r2
  pop r0
endif
pop r2
jmp r2
  
```

| | |
|----|---|
| r2 | 2 |
| r3 | 6 |

← sp (r1)

| | |
|-----------|--|
| loc: cal0 | |
| 2 | |

Stack

```

mult
  psh r2
  loa r2 r1 4
  bgt r2 r0 else
  adc r3 r0 0
  brs endif
else
  sbc r2 r2 1

  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  loa r2 r1 2
  add r3 r3 r2
  pop r0
endif
pop r2
jmp r2
  
```

| | |
|----|---|
| r2 | 2 |
| r3 | 6 |

← sp (r1)

| | |
|-----------|--|
| loc: cal0 | |
| 2 | |

Stack


```

mult
  psh r2
  loa r2 r1 4

  bgt r2 r0 else
  adc r3 r0 0
  brs endif
else
  sbc r2 r2 1

  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  loa r2 r1 2
  add r3 r3 r2
  pop r0
endif
pop r2
jmp r2
    
```

| | |
|----|---|
| r2 | 1 |
| r3 | 6 |

← sp (r1)

| | |
|-----------|---|
| loc: cal0 | 2 |
|-----------|---|

Stack

```

mult
  psh r2
  loa r2 r1 4

  bgt r2 r0 else
  adc r3 r0 0
  brs endif
else
  sbc r2 r2 1

  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  loa r2 r1 2
  add r3 r3 r2
  pop r0
endif
pop r2
jmp r2
    
```

| | |
|----|---|
| r2 | 1 |
| r3 | 6 |

← sp (r1)

| | |
|-----------|---|
| loc: cal0 | 2 |
|-----------|---|

Stack

```

mult
  psh r2
  loa r2 r1 4

  bgt r2 r0 else
  adc r3 r0 0
  brs endif
else
  sbc r2 r2 1

  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  loa r2 r1 2
  add r3 r3 r2
  pop r0
endif
pop r2
jmp r2
    
```

| | |
|----|---|
| r2 | 1 |
| r3 | 6 |

Why psh r3?

← sp (r1)

| | |
|-----------|---|
| 6 | |
| loc: cal0 | 2 |

Stack

```

mult
  psh r2
  loa r2 r1 4

  bgt r2 r0 else
  adc r3 r0 0
  brs endif
else
  sbc r2 r2 1

  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  loa r2 r1 2
  add r3 r3 r2
  pop r0
endif
pop r2
jmp r2
    
```

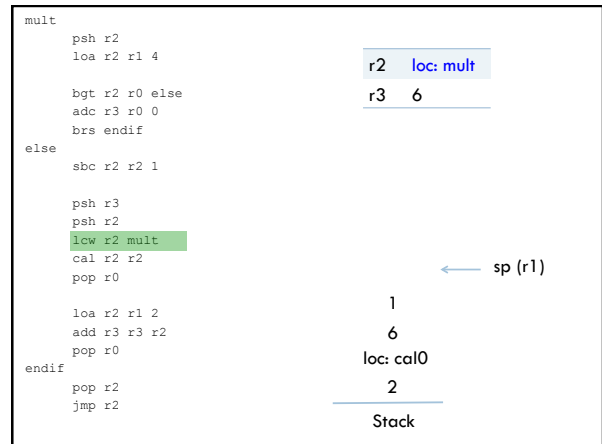
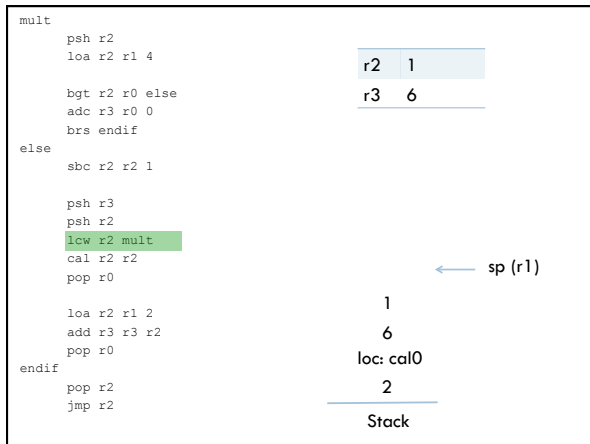
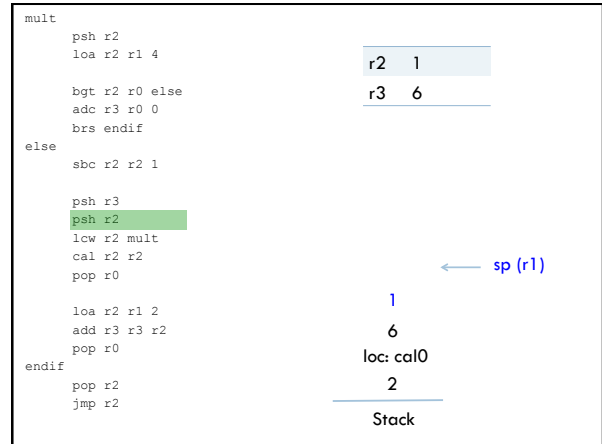
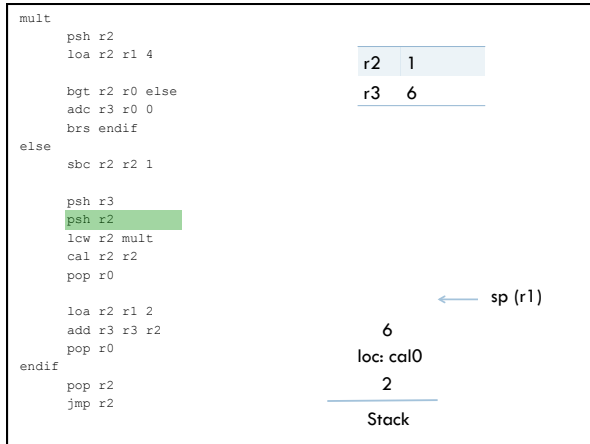
| | |
|----|---|
| r2 | 1 |
| r3 | 6 |

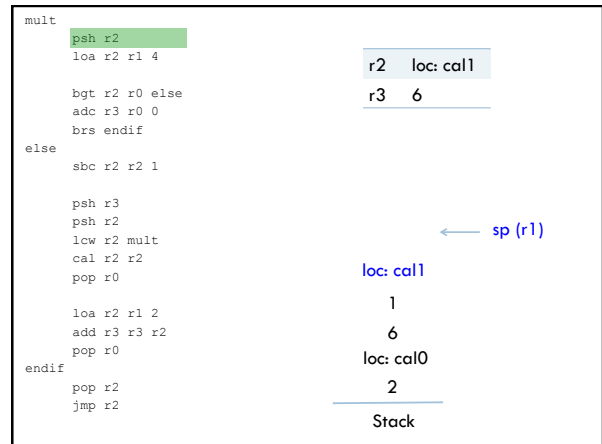
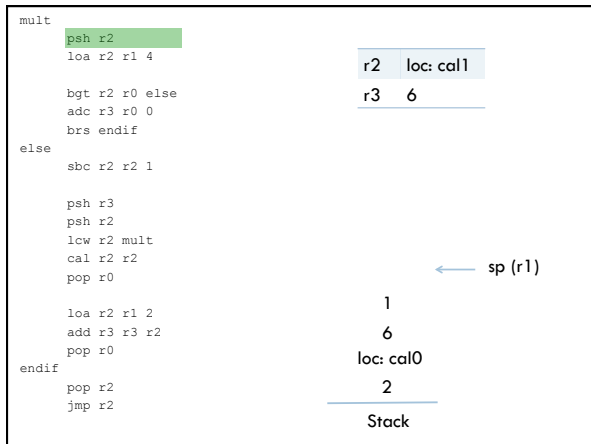
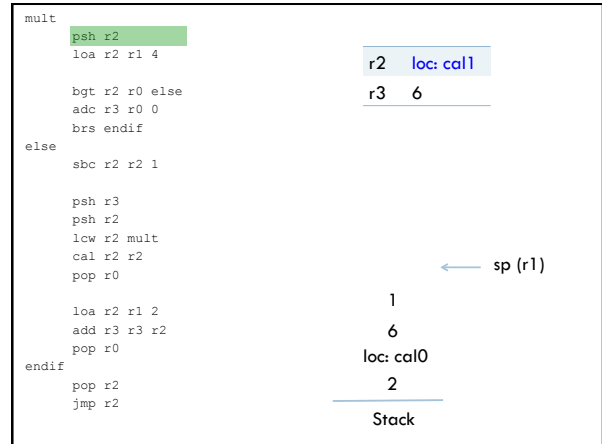
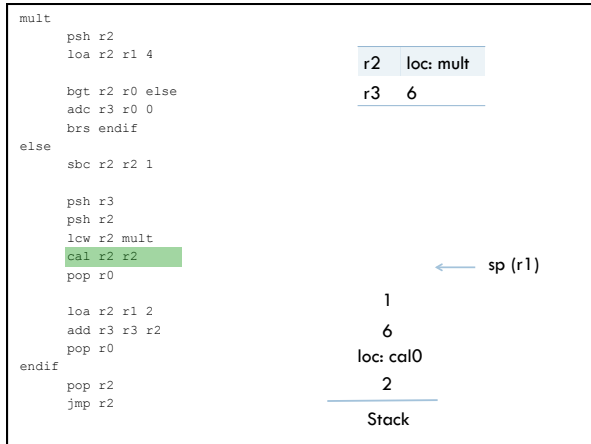
← sp (r1)

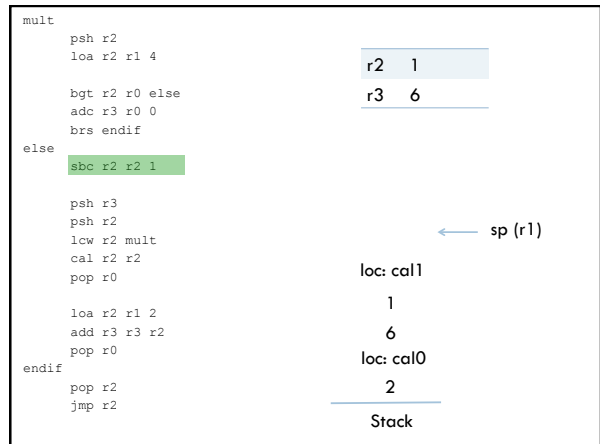
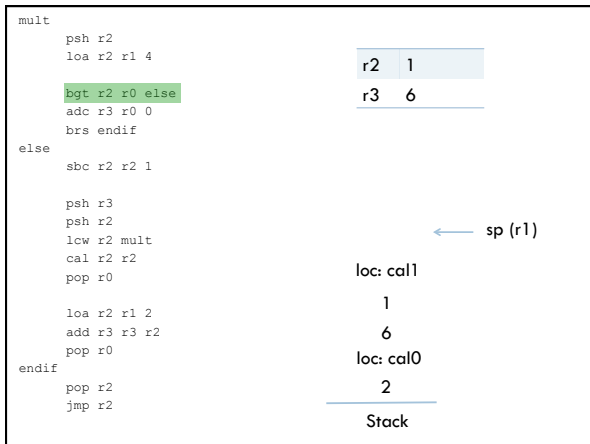
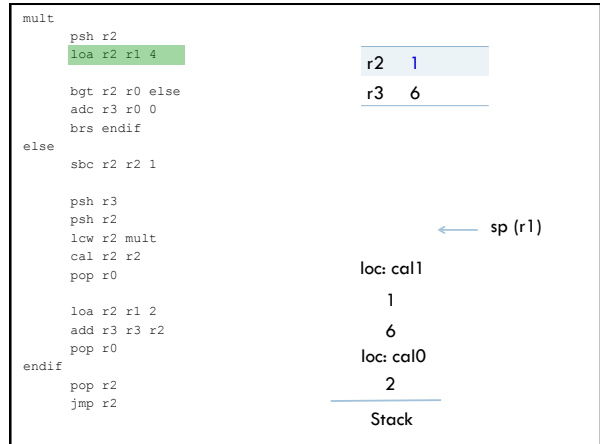
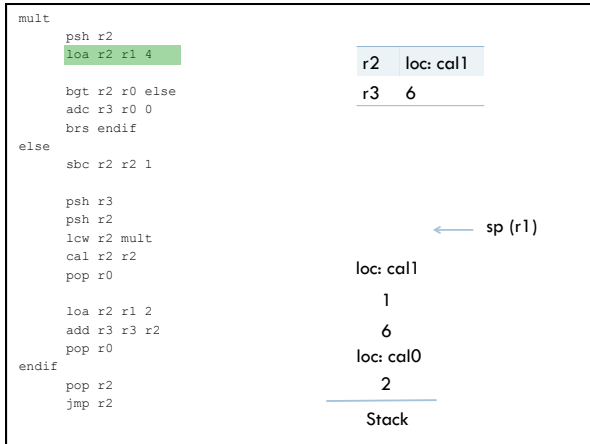
| | |
|-----------|---|
| 6 | |
| loc: cal0 | 2 |

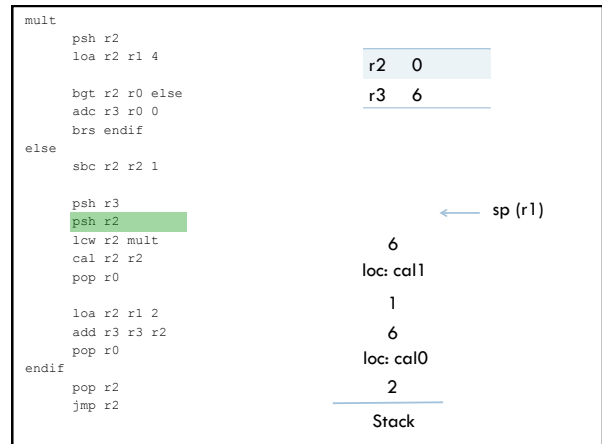
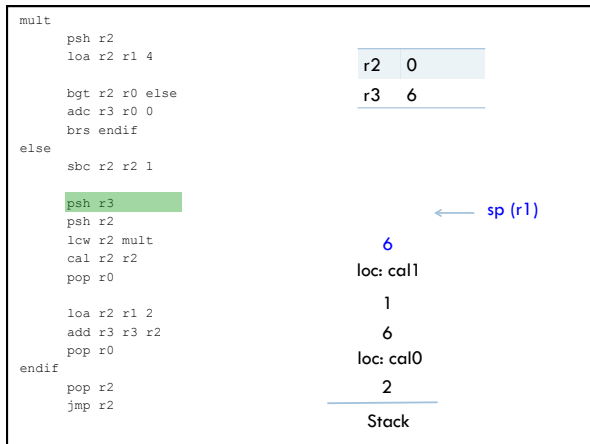
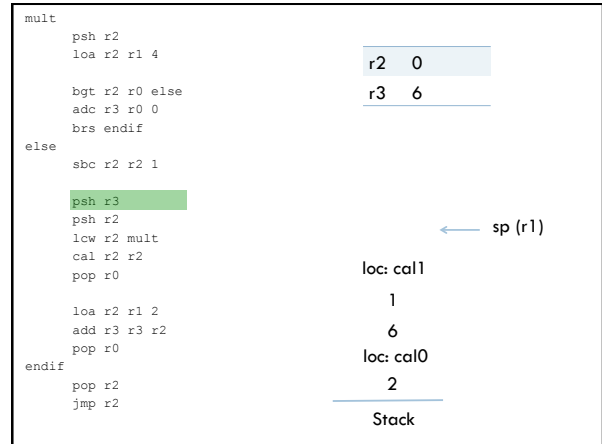
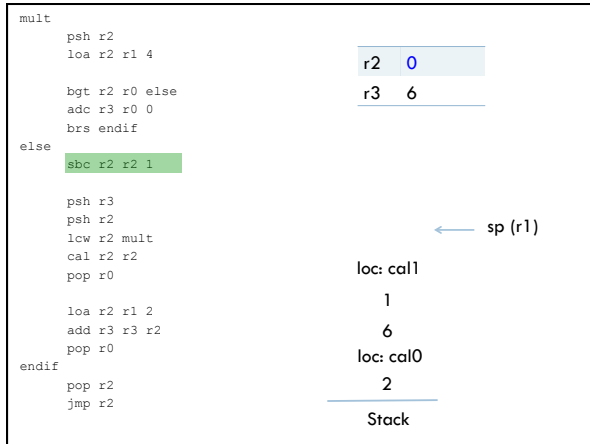
Stack

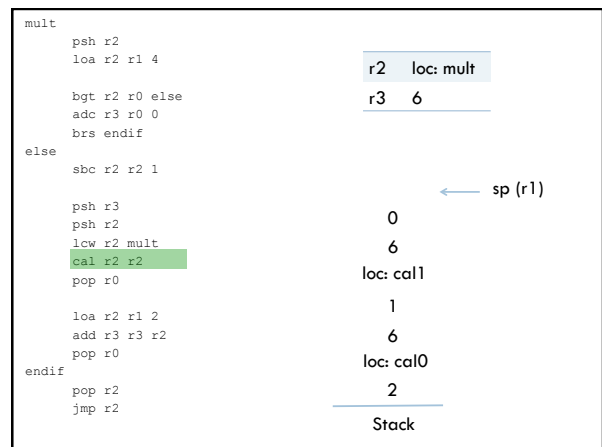
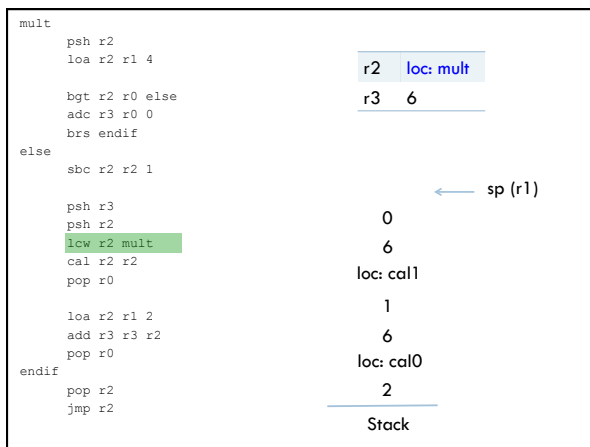
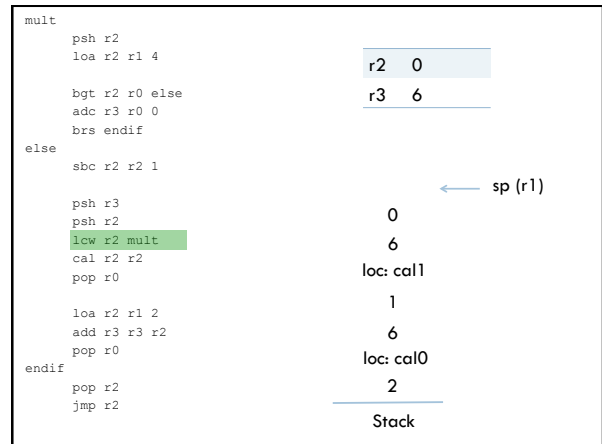
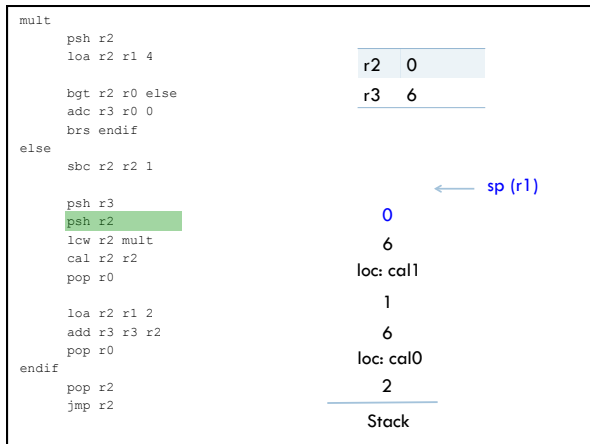
- We're about to make a function call
- The result of that call will go into r3 so we'll lose what's in there if we don't save it!

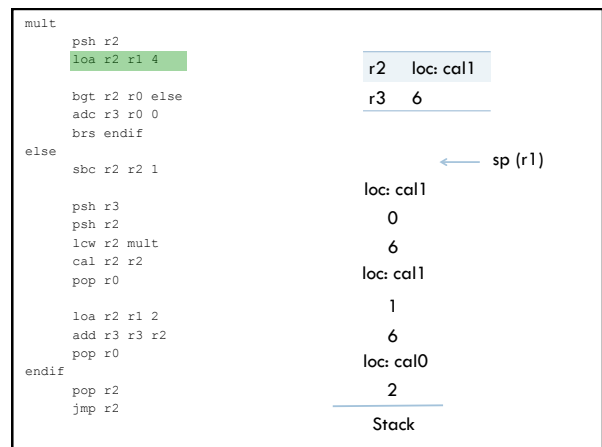
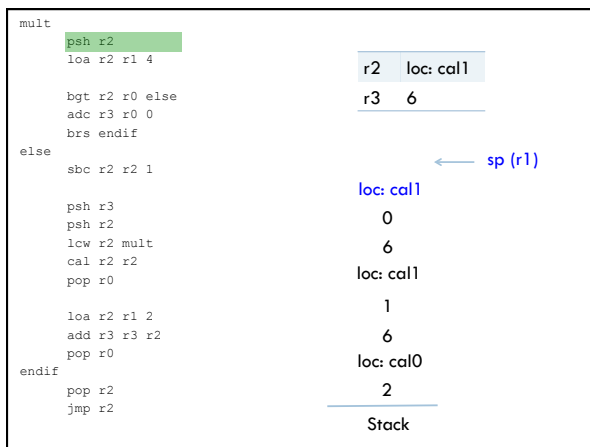
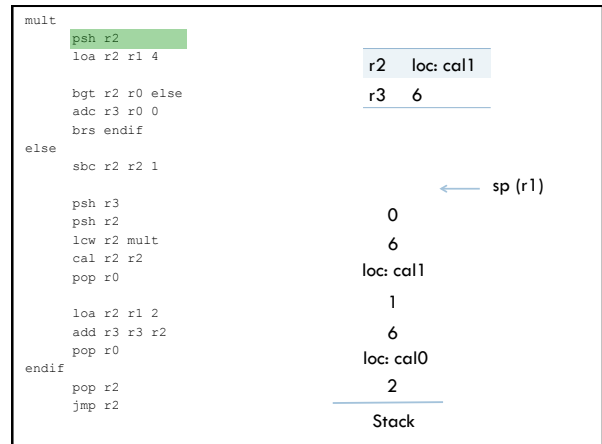
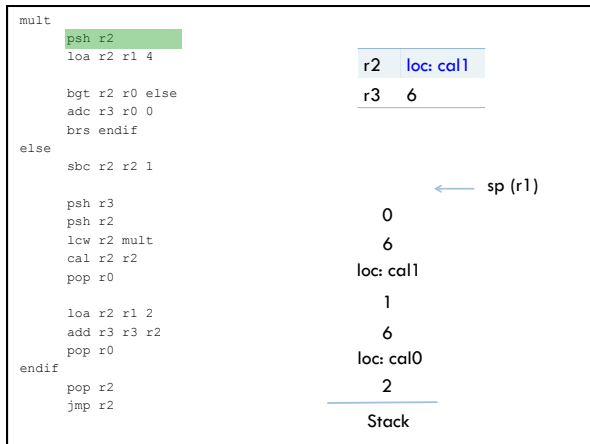


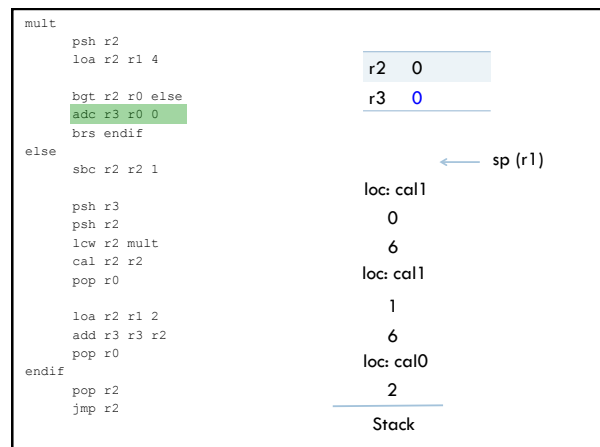
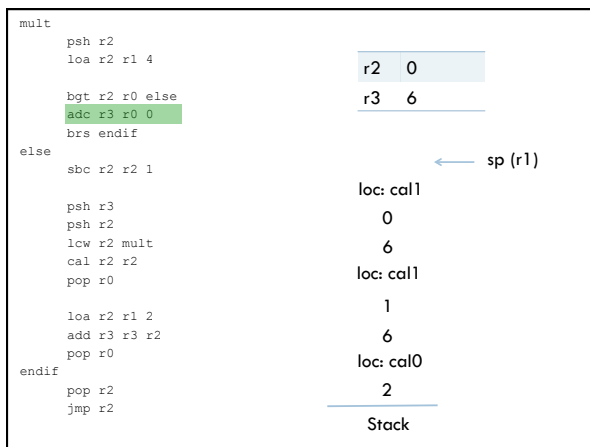
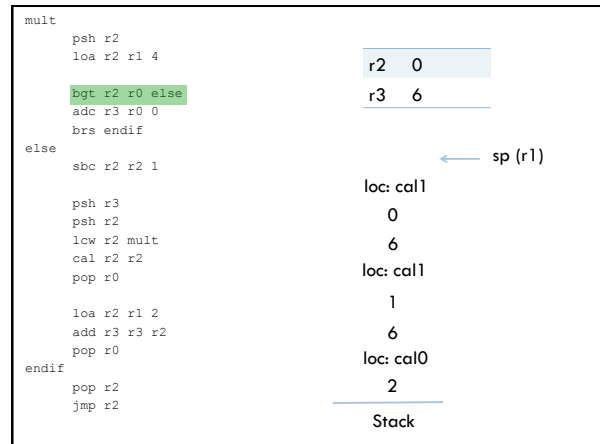
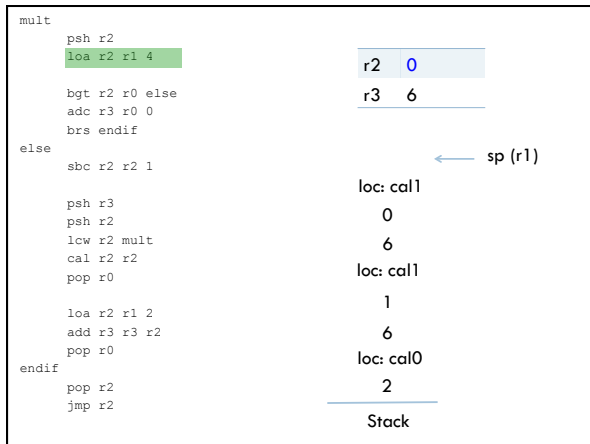


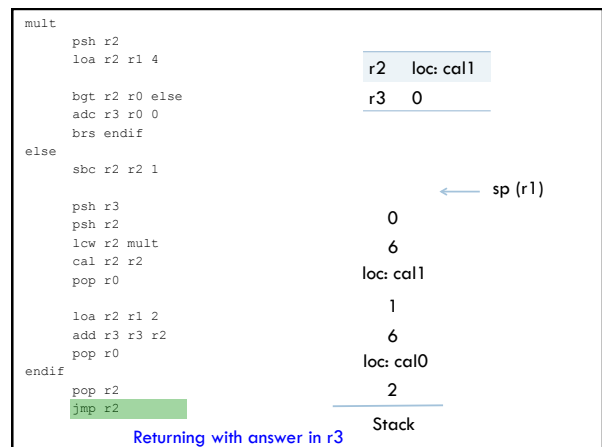
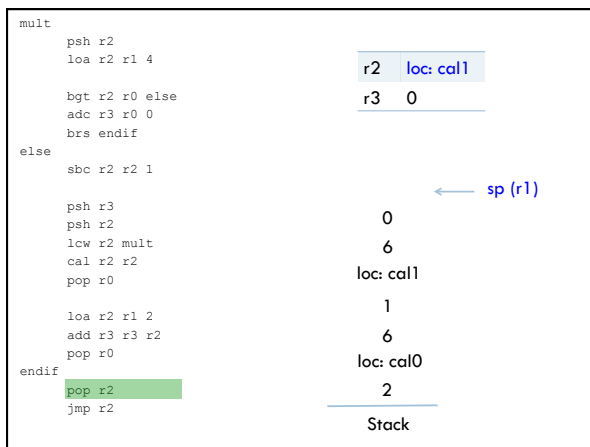
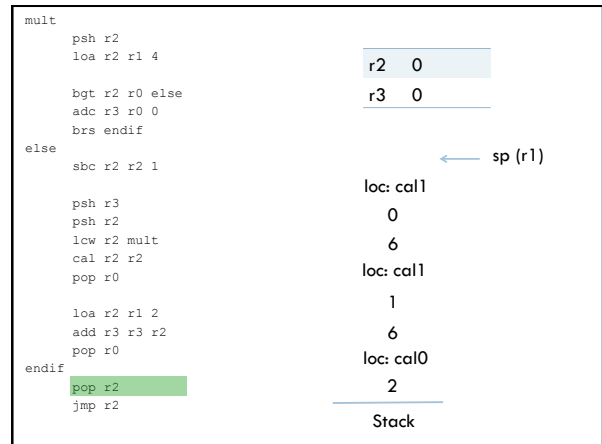
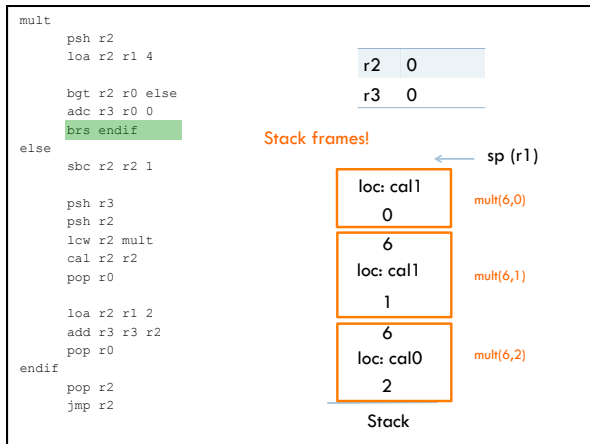


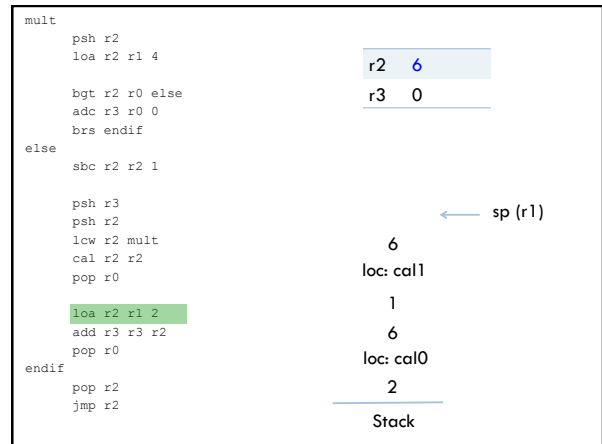
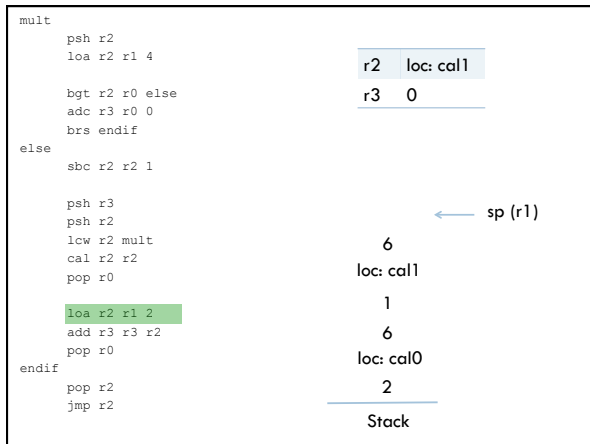
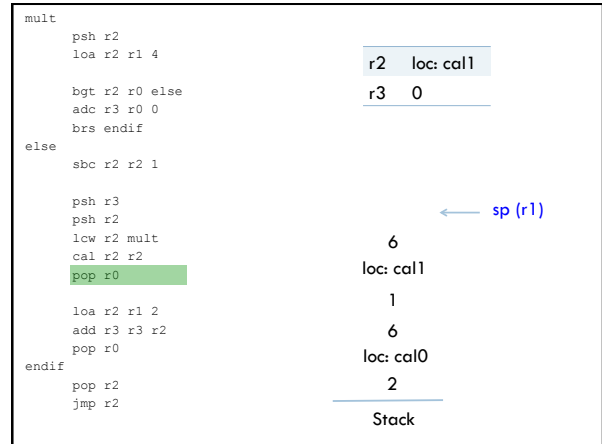
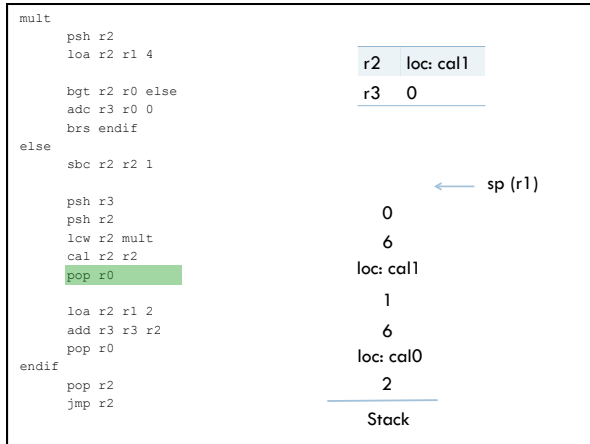


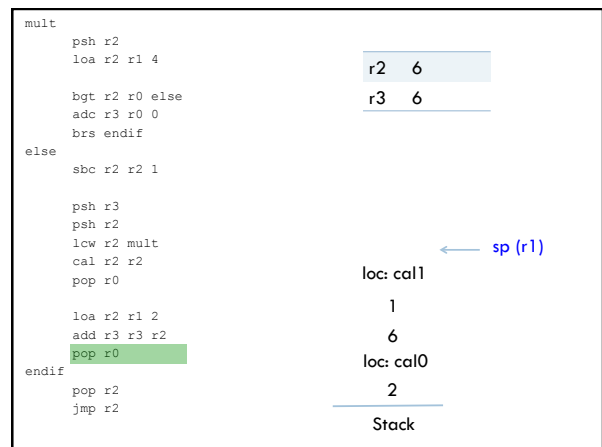
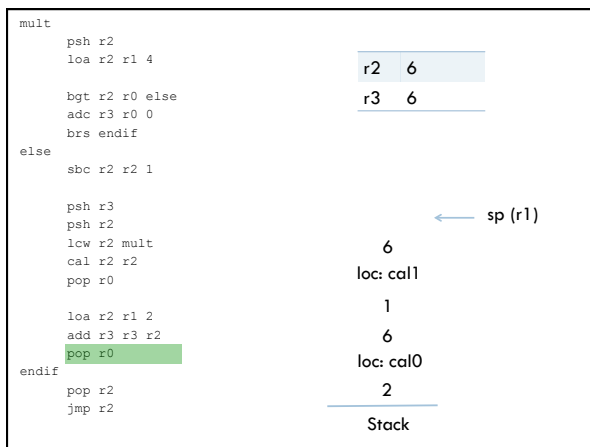
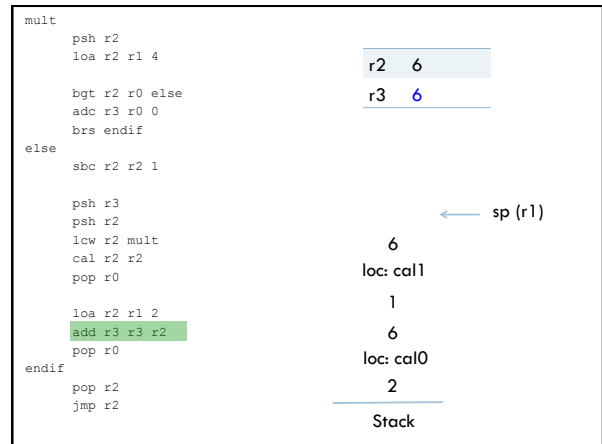
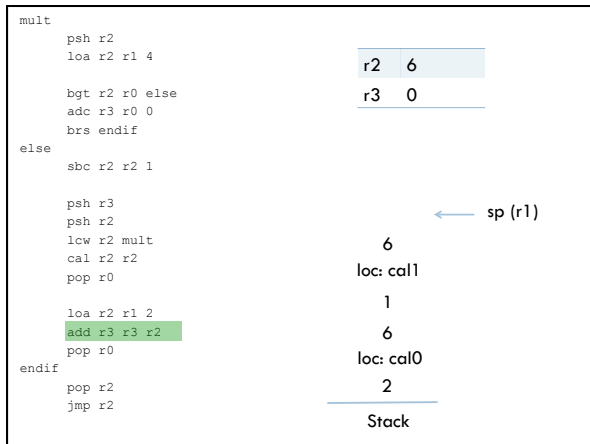


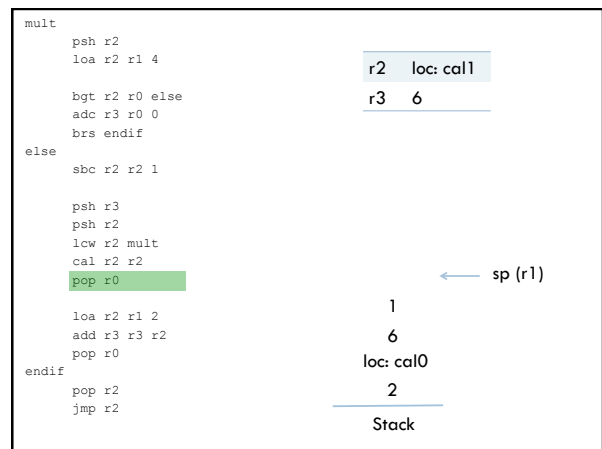
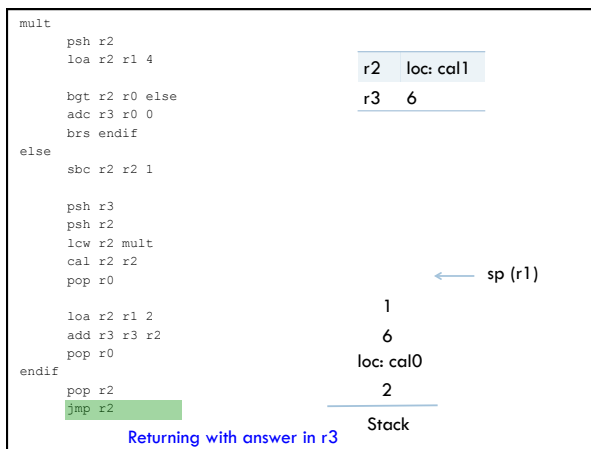
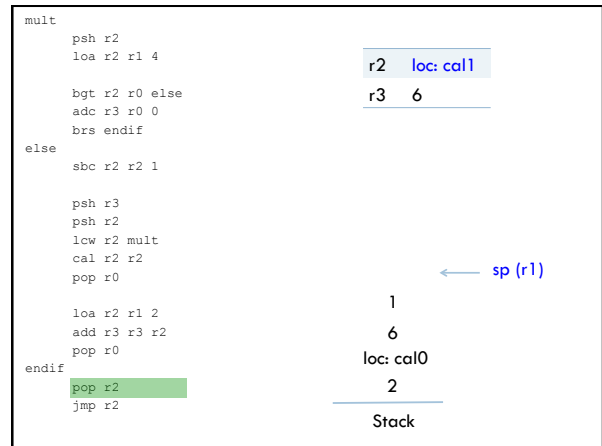
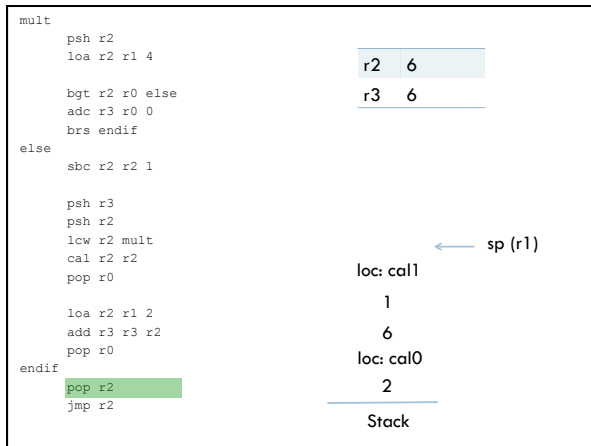


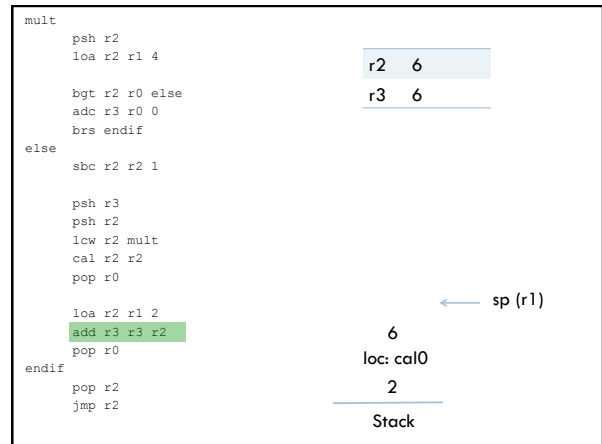
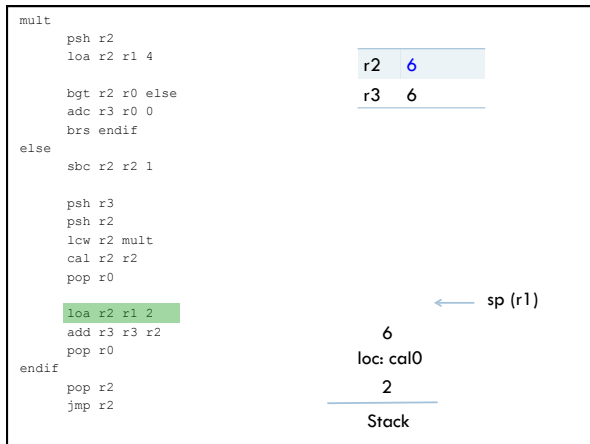
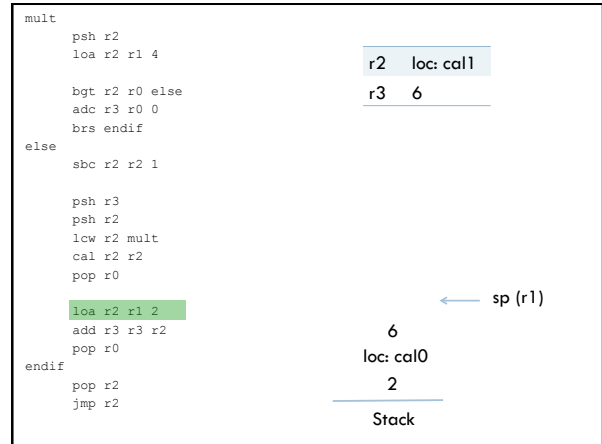
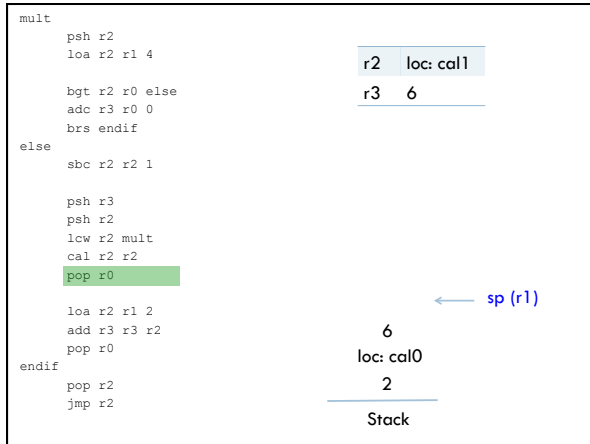


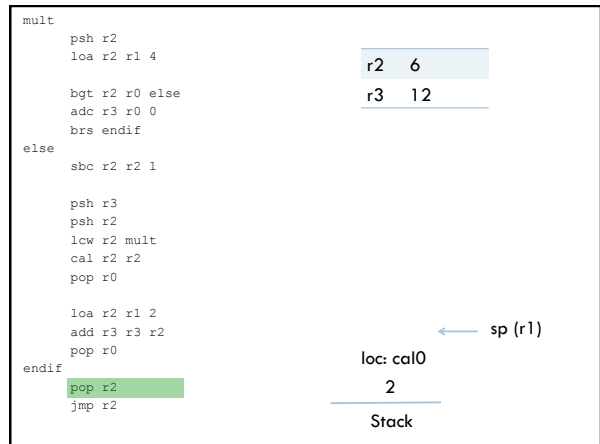
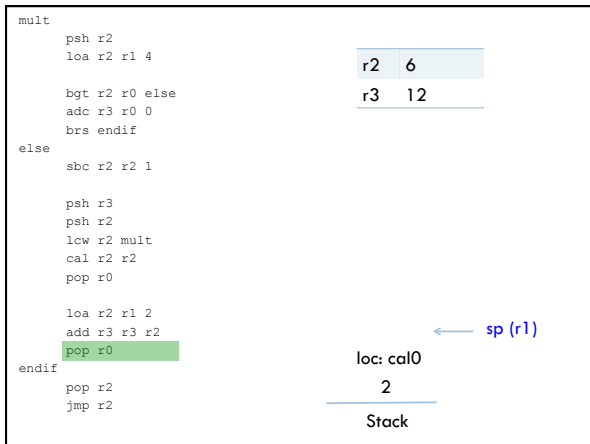
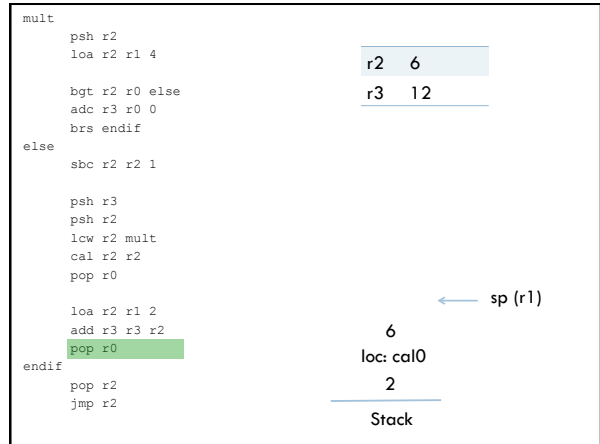
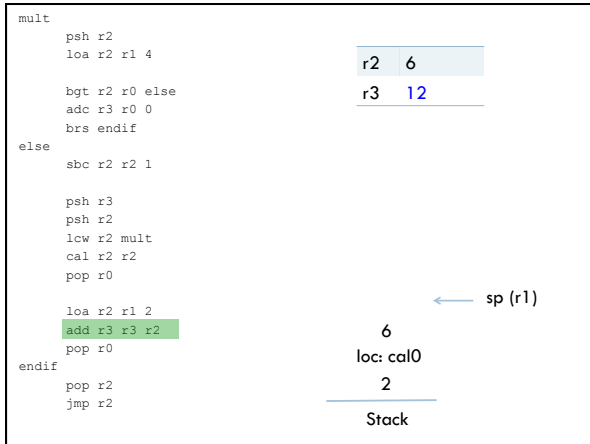


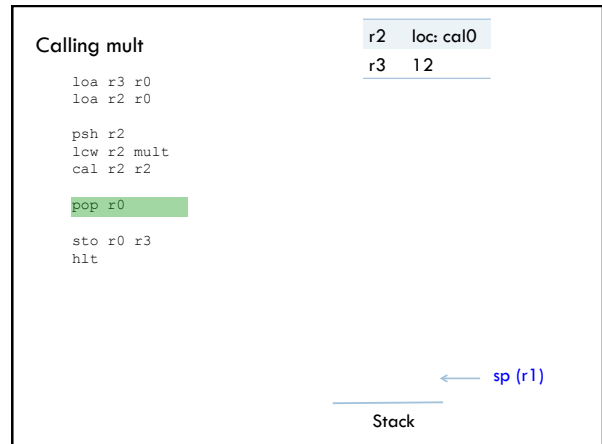
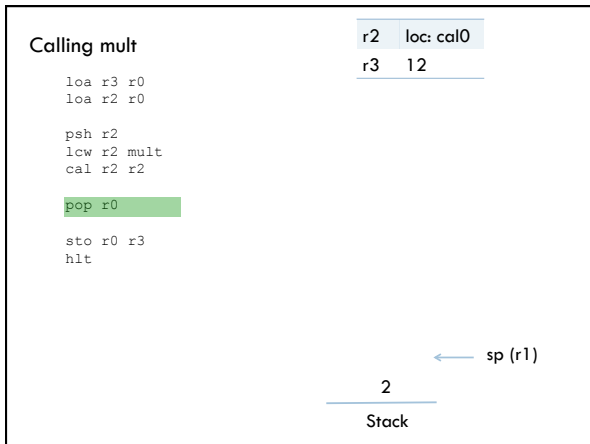
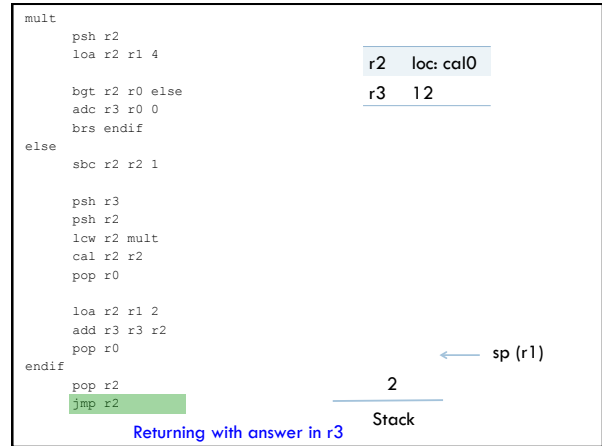
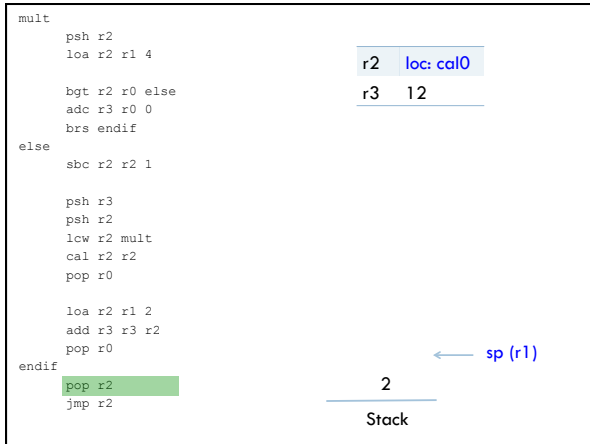


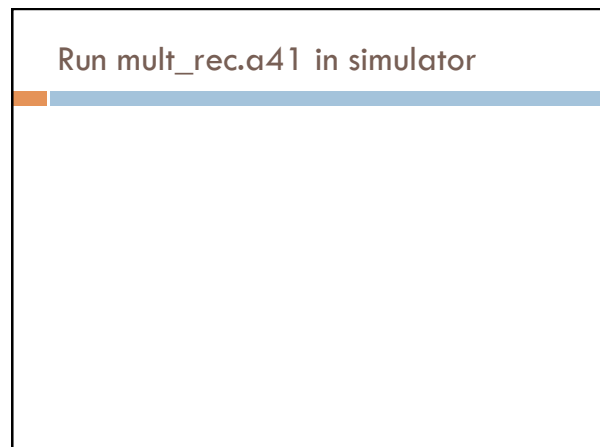
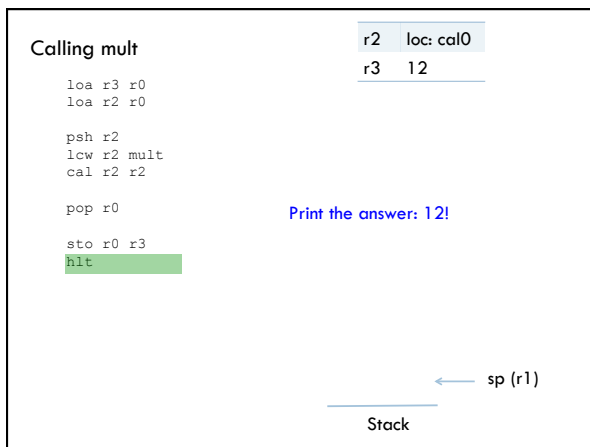
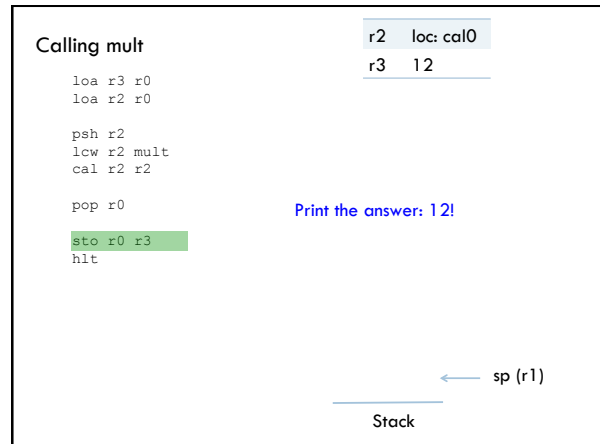
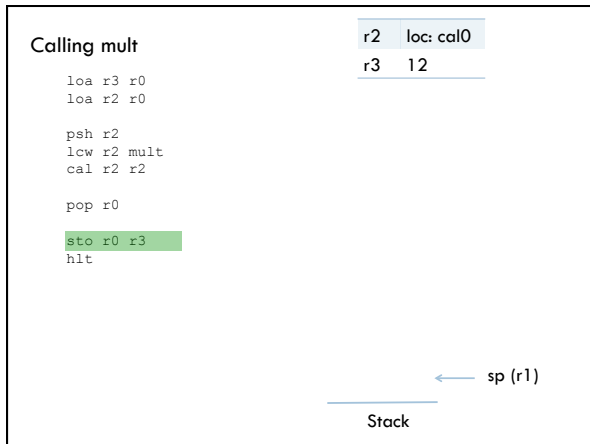












CS41B programming advice

1. Match your psh and pops
2. Follow the register conventions
3. Develop code incrementally
4. Debugging: write out stack, registers, etc. on paper and compare against system execution

Examples from this lecture

<http://www.cs.pomona.edu/~dkauchak/classes/cs52/examples/cs41b/>