

Sorting Conclusions

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cs302
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Administrative

- Find a partner to bounce ideas off of
- Assignment 2
 - Proofs should be very concrete
 - Proving big-O, must satisfy definition
 - When comparing (i.e. ordering) functions you're not sure about, set them equal to each other and do some manipulations
 - e.g. compare n and $3^{\sqrt{\log n}}$ or n^n and 3^{3^n}
 - Be careful of things that seem "too good to be true", e.g. a new $O(n \log n)$ sorting algorithm that you've never heard of ☺



Administrivia continued

- Assignment 3
 - Substitution method is a proof by induction. Follow the steps we talked about before.
 - Make sure to prove/justify any non-trivial steps (e.g. showing big-O, Ω and θ)
 - If you're stuck on an algorithm, brainstorm different possibilities
 - Even start with the naive algorithm and think about where it's inefficient
- Latex/writeups
 - Use vertical space to help make your argument/steps more clear
 - Look at what is being generated and make sure it looks like you expect (e.g. $n^{2.5}$ vs. $n^{\{2.5\}}$ or $n^{1+\epsilon}$ vs $n^{\{1 + \epsilon\}}$)



Sorting bounds

Mergesort is $O(n \log n)$

Quicksort is $O(n \log n)$ on average

Can we do better?



Comparison-based sorting

Sorted order is determined based **only** on a comparison between input elements

- $A[i] < A[j]$
- $A[i] > A[j]$
- $A[i] = A[j]$
- $A[i] \leq A[j]$
- $A[i] \geq A[j]$

Do any of the sorting algorithms we've looked at use additional information?

- No
- All the algorithms we've seen are comparison-based sorting algorithms



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In Java (and many languages) for a class of objects to be sorted we define a comparator

What does it do?



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In Java (and many languages) for a class of objects to be sorted we define a comparator

What does it do?

- Just compares any two elements
- Useful for comparison-based sorting algorithms



Comparison-based sorting

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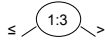
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Can we do better than $O(n \log n)$ for comparison based sorting approaches?

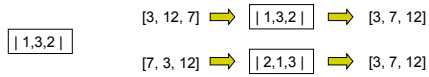


Decision-tree model

- Full binary tree representing the comparisons between elements by a sorting algorithm
- Internal nodes contain indices to be compared

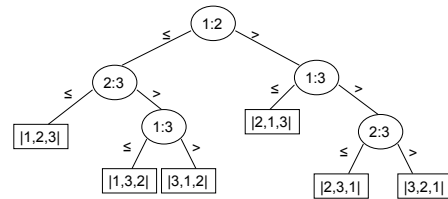


- Leaves contain a complete permutation of the input

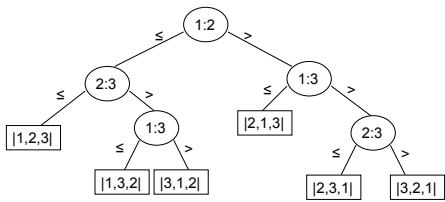


- Tracing a path from root to leaf gives the correct reordering/permutation of the input for an input

A decision tree model

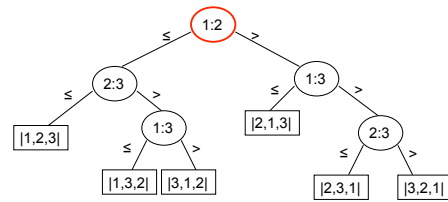


A decision tree model



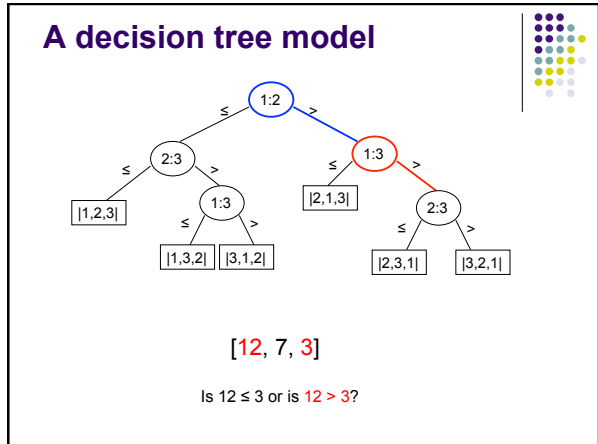
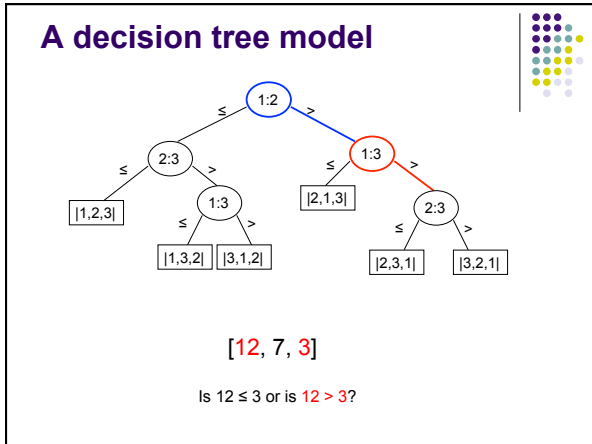
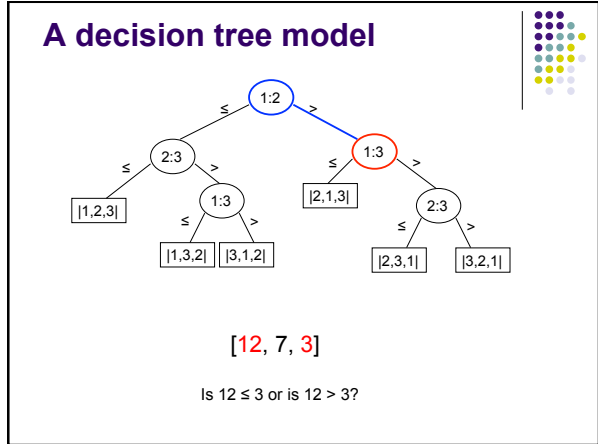
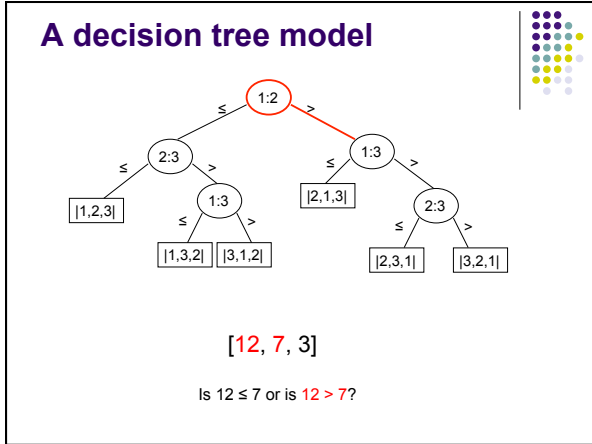
[12, 7, 3]

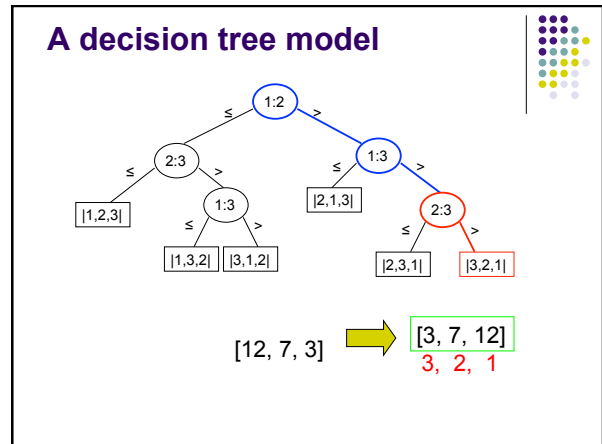
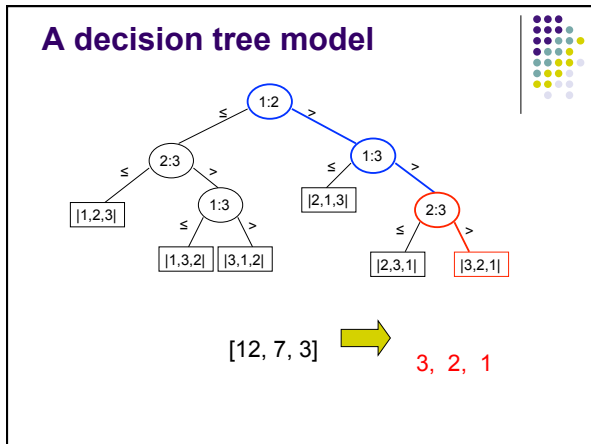
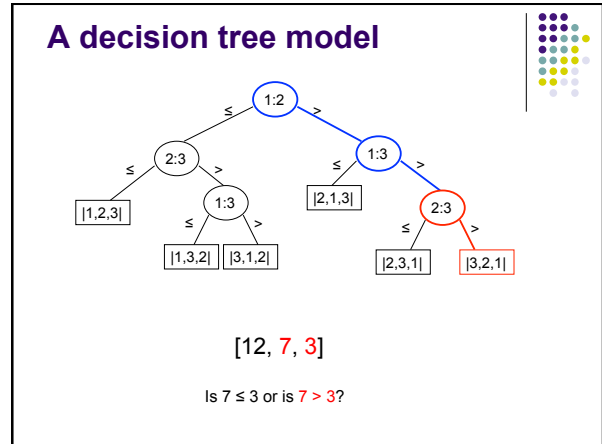
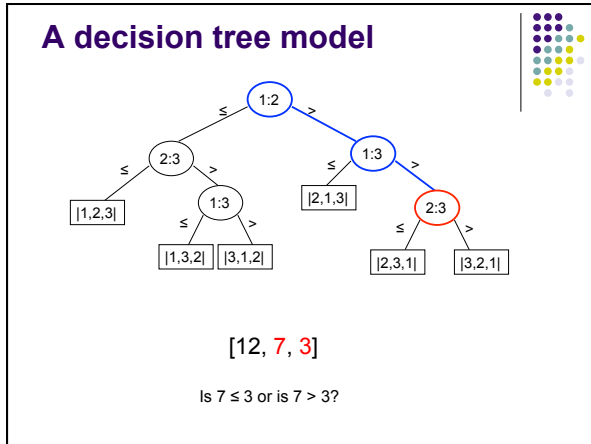
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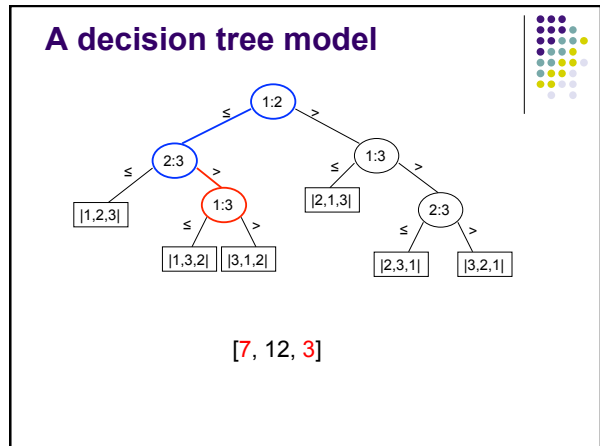
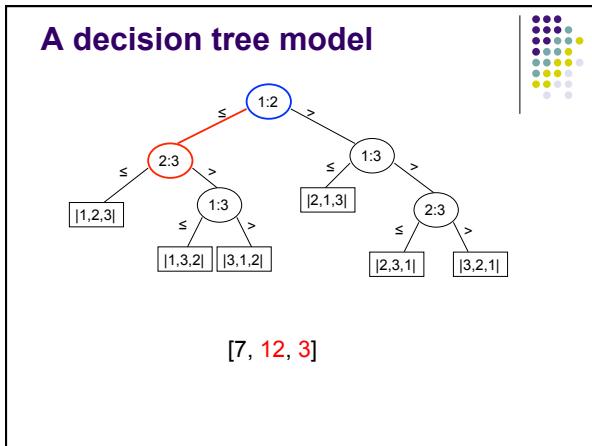
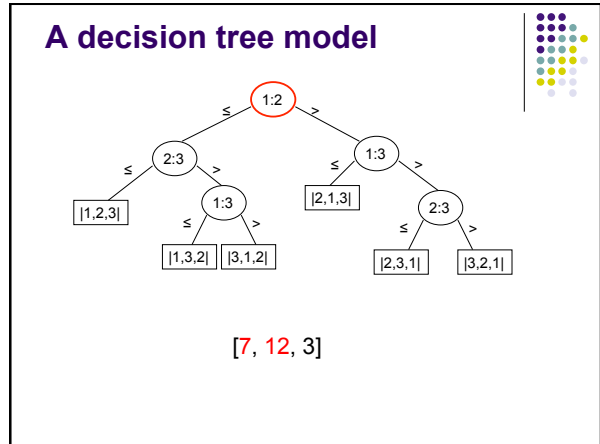
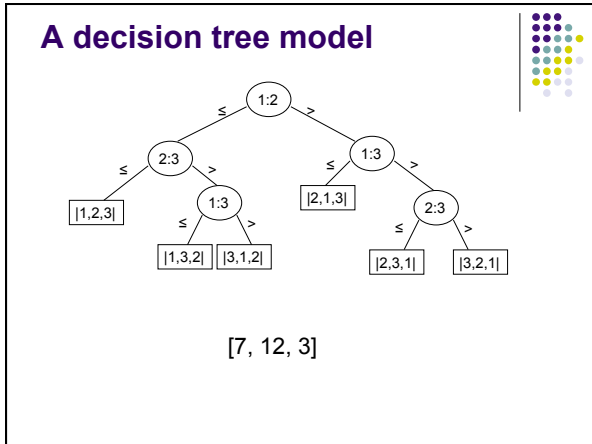


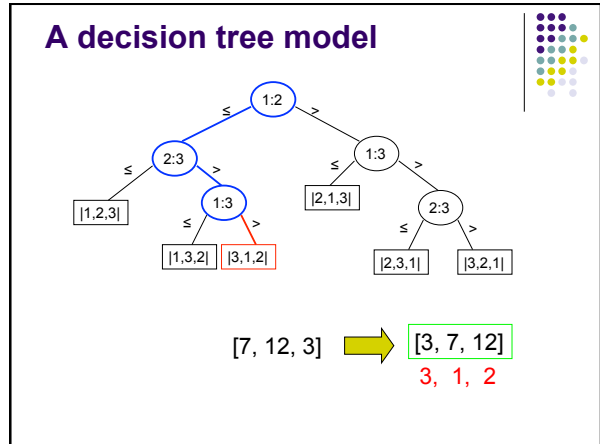
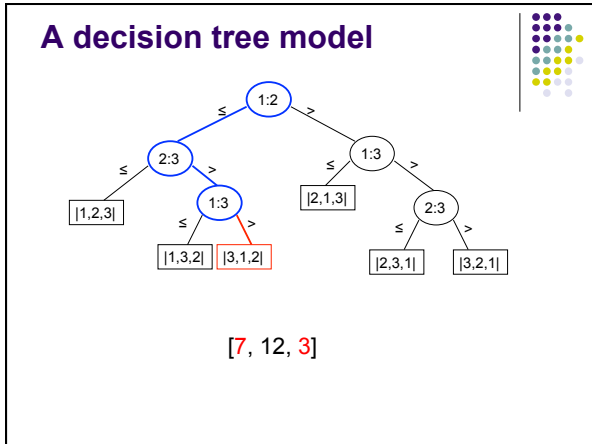
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Is $12 \leq 7$ or is $12 > 7$?









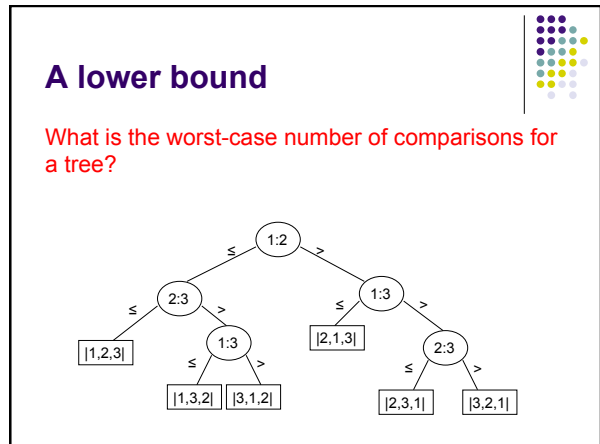
How many leaves are in a decision tree?

Leaves **must** have all possible permutations of the input

What if decision tree model didn't?

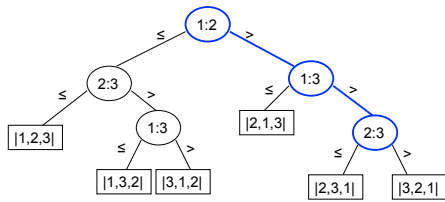
Some input would exist that didn't have a correct reordering

Input of size n , $n!$ leaves



A lower bound

The longest path in the tree, i.e. the height



A lower bound

What is the maximum number of leaves a binary tree of height h can have?

A complete binary tree has 2^h leaves

$$2^h \geq n!$$

$$h \geq \log n!$$

$$h = \Omega(n \log n) \quad \text{from hw } \odot$$

Can we do better?