NP-COMPLETE REDUCTIONS

CS302, Spring 2013
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Admin

- Last assignment out today (yay!)
- Review topics?
  - E-mail me if you have others…
- CS senior theses
  - Tue 3-4, Wed 3-4:30, Thur 3-4 in MBH 104

NP problems

NP is the set of problems that can be *verified* in polynomial time

A problem can be verified in polynomial time if you can check that a given solution is correct in polynomial time

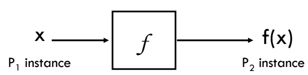
(NP is an abbreviation for non-deterministic polynomial time)

Reduction function

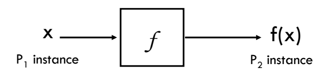
Reduction function

Given two problems P_1 and P_2 *a reduction function*, $f(x)$, is a function that transforms a problem instance x of type P_1 to a problem instance of type P_2

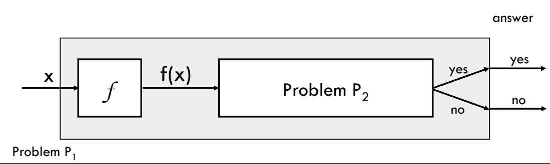
such that: a solution to x exists for P_1 iff a solution for $f(x)$ exists for P_2



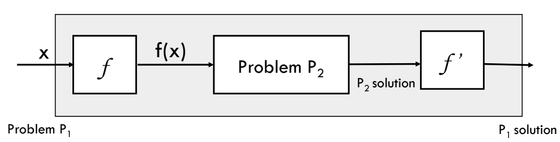
Reduction function



Allows us to solve P_1 problems if we have a solver for P_2

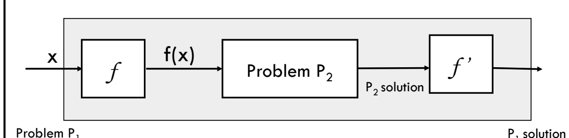


Reduction function



Most of the time we'll worry about yes no question, however, if we have more complicated answers we often just have to do a little work to the solution to the problem of P_2 to get the answer

Reduction function: Example

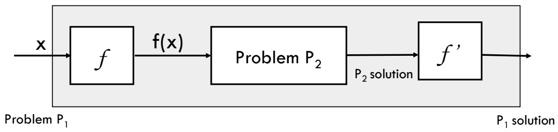


P1 = Bipartite matching
P2 = Network flow

Reduction function (f): Given *any* bipartite matching problem turn it into a network flow problem

What is f and what is f'?

Reduction function: Example



P1 = Bipartite matching
P2 = Network flow

Reduction function (f): Given *any* bipartite matching problem turn it into a network flow problem

A reduction function reduces problems instances

NP-Complete

A problem is *NP-complete* if:

1. it can be verified in polynomial time (i.e. in NP)
2. *any* NP-complete problem can be reduced to the problem in polynomial time (is NP-hard)

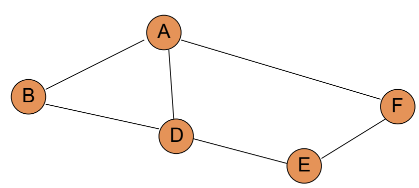
Why are NP-complete problems interesting?

NP-Complete problems

What are some of the NP-complete problems we've talked about (or that you know about)?

Hamiltonian cycle

Given an undirected graph G=(V, E), a hamiltonian cycle is a cycle that visits every vertex V exactly once



Longest path

Given a graph G with nonnegative edge weights does a simple path exist from *s* to *t* with weight at least *g*?

3D matching

Bipartite matching: given two sets of things and pair constraints, find a matching between the sets

3D matching: given three sets of things and triplet constraints, find a matching between the sets

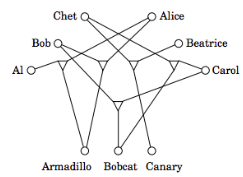


Figure from Dasgupta et. al 2008

3-SAT

A boolean formula is in *n-conjunctive normal form* (*n*-CNF) if:

- it is expressed as an AND of clauses
- where each clause is an OR of no more than *n* variables

$$(a \vee \neg a \vee \neg b) \wedge (c \vee b \vee d) \wedge (\neg a \vee \neg c \vee \neg d)$$

3-SAT: Given a 3-CNF boolean formula, is it satisfiable?

SAT

Given a boolean formula of *n* boolean variables joined by *m* connectives (AND, OR or NOT) is there a setting of the variables such that the boolean formula evaluate to true?

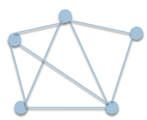
$$(a \wedge b) \vee (\neg a \wedge \neg b)$$

$$((\neg(b \vee \neg c) \wedge a) \vee (a \wedge b \wedge c)) \wedge c \wedge \neg b$$

CLIQUE

A *clique* in an undirected graph G = (V, E) is a subset V' ⊆ V of vertices that are fully connected, i.e. every vertex in V' is connected to every other vertex in V'

CLIQUE problem: Does G contain a clique of size k?



Is there a clique of size 4 in this graph?

Proving NP-completeness

Given a problem NEW to show it is NP-Complete

1. Show that NEW is in NP
   a. Provide a verifier
   b. Show that the verifier runs in polynomial time
2. Show that all NP-complete problems are reducible to NEW in polynomial time
   a. Describe a reduction function *f* from a known NP-Complete problem to NEW
   b. Show that *f* runs in polynomial time
   c. Show that a solution exists to the NP-Complete problem IFF a solution exists *to the NEW problem generate by f*

Proving NP-completeness

Show that a solution exists to the NP-Complete problem IFF a solution exists *to the NEW problem generate by f*

- Assume we have an NP-Complete problem instance that has a solution, show that the NEW problem instance generated by *f* has a solution
- Assume we have a problem instance of NEW *generated by f* that has a solution, show that we can derive a solution to the NP-Complete problem instance

Other ways of proving the IFF, but this is often the easiest

HALF-CLIQUE

Given a graph G, does the graph contain a clique containing exactly half the vertices?

Is HALF-CLIQUE an NP-complete problem?

Is Half-Clique NP-Complete?

1. Show that NEW is in NP
   a. Provide a verifier
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Given a graph G, does the graph contain a clique containing exactly half the vertices?

HALF-CLIQUE

1. Show that HALF-CLIQUE is in NP
   a. Provide a verifier
   b. Show that the verifier runs in polynomial time

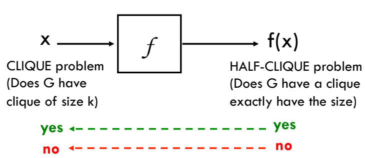
Verifier: A solution consists of the set of vertices in V'
- check that |V '| = |V|/2
- for all pairs of u, v ∈ V'
  - there exists an edge (u,v) ∈ E

- Check for edge existence in O(V)
- $O(V^2)$ checks
- $O(V^3)$ overall, which is polynomial

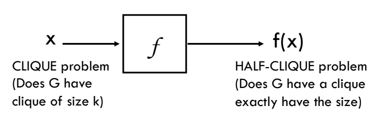
HALF-CLIQUE

2. Show that all NP-complete problems are reducible to HALF-CLIQUE in polynomial time
   a. Describe a reduction function *f* from a known NP-Complete problem to HALF-CLIQUE
   b. Show that *f* runs in polynomial time
   c. Show that a solution exists to the NP-Complete problem IFF a solution exists *to the HALF-CLIQUE problem generate by f*

Reduce CLIQUE to HALF-CLIQUE:
Given a problem instance of CLIQUE, turn it into a problem instance of HALF-CLIQUE



HALF-CLIQUE



Three cases:
1. k = |V|/2

2. k < |V|/2

3. k > |V|/2

HALF-CLIQUE

Reduce CLIQUE to HALF-CLIQUE:
Given an instance of CLIQUE, turn it into an instance of HALF-CLIQUE

It's already a half-clique problem

$f(G, k)$
if $\lceil |V| \rceil / 2 = k$
    return G
elseif $k < \lceil |V| \rceil / 2$
    return G plus $(|V| - 2k)$ nodes which are fully connected and are connected to every node in V
else
    return G plus $2k - |V|$ nodes which have no edges

HALF-CLIQUE

Reduce CLIQUE to HALF-CLIQUE:
Given an instance of CLIQUE, turn it into an instance of HALF-CLIQUE

We're looking for a clique that is smaller than half, so add an artificial clique to the graph and connect it up to all vertices

$f(G, k)$
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HALF-CLIQUE

Reduce CLIQUE to HALF-CLIQUE:
Given an instance of CLIQUE, turn it into an instance of HALF-CLIQUE

We're looking for a clique that is bigger than half, so add vertices until k = |V|/2

$f(G, k)$
if $\lceil |V| \rceil / 2 = k$
    return G
elseif $k < \lceil |V| \rceil / 2$
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    return G plus $2k - |V|$ nodes which have no edges

Runtime: From the construction we can see that it is polynomial time

Reduction proof

Show that a solution exists to the NP-Complete problem IFF a solution exists *to the NEW problem generate by f*

- Assume we have an NP-Complete problem instance that has a solution, show that the NEW problem instance generated by *f* has a solution
- Assume we have a problem instance of NEW *generated by f* that has a solution, show that we can derive a solution to the NP-Complete problem instance

```
f(G,k)
if ⌈|V|⌉/2 = k
        return G
elseif k < ⌈|V|⌉/2
        return G plus (|V| − 2k) nodes which are fully connected
        and are connected to every node in V
else
        return G plus 2k − |V| nodes which have no edges
```

Reduction proof

Given a graph G that has a CLIQUE of size k, show that f(G,k) has a solution to HALF-CLIQUE

If k = |V|/2:

- the graph is unmodified
- f(G,k) has a clique that is half the size

Reduction proof

Given a graph G that has a CLIQUE of size k, show that f(G,k) has a solution to HALF-CLIQUE

If k < |V|/2:

- we added a clique of |V|- 2k fully connected nodes
- there are |V| + |V| - 2k = 2(|V|-k) nodes in f(G)
- there is a clique in the original graph of size k
- plus our added clique of |V|-2k
- k + |V|-2k = |V|-k, which is half the size of f(G)

Reduction proof

Given a graph G that has a CLIQUE of size k, show that f(G,k) has a solution to HALF-CLIQUE

If k >|V|/2:

- we added 2k - |V| unconnected vertices
- f(G) contains |V| + 2k - |V| = 2k vertices
- Since the original graph had a clique of size k vertices, the new graph will have a half-clique

Reduction proof

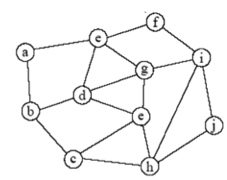
Given a graph f(G) that has a CLIQUE half the elements, show that G has a clique of size k

Key: f(G) was constructed by your reduction function

Use a similar argument to what we used in the other direction

Independent-Set

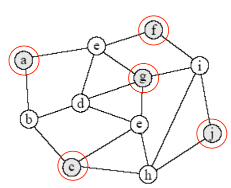
Given a graph G = (V, E) is there a subset V'⊆ V of vertices of size |V '| = k that are independent, i.e. for any pair of vertices u, v ∈ V' there exists no edge between any of these vertices



Does the graph contain an independent set of size 5?

Independent-Set

Given a graph G = (V, E) is there a subset V'⊆ V of vertices of size |V '| = k that are independent, i.e. for any pair of vertices u, v ∈ V' there exists no edge between any of these vertices

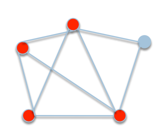


Independent-Set is NP-Complete

CLIQUE revisited

A *clique* in an undirected graph G = (V, E) is a subset V' ⊆ V of vertices that are fully connected, i.e. every vertex in V' is connected to every other vertex in V'

CLIQUE problem: Does G contain a clique of size k?



Is CLIQUE NP-Complete?

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   a. Provide a verifier
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   a. Describe a reduction function *f* from a known NP-Complete problem to CLIQUE
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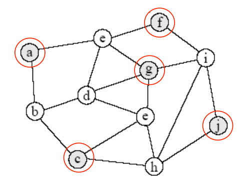
   c. Show that a solution exists to the NP-Complete problem IFF a solution exists *to the CLIQUE problem generate by f*

Given a graph G, does the graph contain a clique containing exactly half the vertices?

Independent-Set

Given a graph G = (V, E) is there a subset V'⊆ V of vertices of size |V '| = k that are independent, i.e. for any pair of vertices u, v ∈ V' there exists no edge between any of these vertices. Is there an independent set of size k?



Reduce Independent-Set to CLIQUE

Independent-Set to Clique

Given a graph G = (V, E) is there a subset V'⊆ V of vertices of size |V '| = k that are independent, i.e. for any pair of vertices u, v ∈ V' there exists no edge between any of these vertices

Both are selecting vertices

Independent set wants vertices where NONE are connected

Clique wants vertices where ALL are connected

How can we convert a NONE problem to an ALL problem?

Independent-Set to Clique

Given a graph G = (V, E), the complement of that graph G' = (V, E) is the a graph constructed by remove all edges E and including all edges not in E

For example, for adjacency matrix this is flipping all of the bits

```
f(G)
    return G'
```

Reduction proof

Show that a solution exists to the NP-Complete problem IFF a solution exists *to the NEW problem generate by f*

- Assume we have an Independent-Set problem instance that has a solution, show that the Clique problem instance generated by *f* has a solution
- Assume we have a problem instance of Clique *generated by f* that has a solution, show that we can derive a solution to Independent-Set problem instance

f(G)
    return G'

Proof

Given a graph G that has an independent set of size k, show that f(G) has a clique of size k

- By definition, the independent set has no edges between any vertices
- These will all be edges in f(G) and therefore they will form a clique of size k

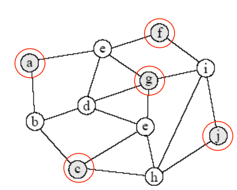
Proof

Given f(G) that has clique of size k, show that G has an independent set of size k

- By definition, the clique will have an edge between every vertex
- None of these vertices will therefore be connected in G, so we have an independent set

Independent-Set revisited

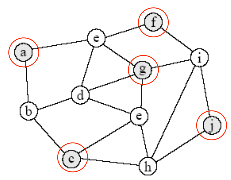
Given a graph G = (V, E) is there a subset V'⊆ V of vertices of size |V '| = k that are independent, i.e. for any pair of vertices u, v ∈ V' there exists no edge between any of these vertices



Is Independent-Set NP-Complete?

Independent-Set revisited

Given a graph G = (V, E) is there a subset V'⊆ V of vertices of size |V '| = k that are independent, i.e. for any pair of vertices u, v ∈ V' there exists no edge between any of these vertices



Reduce 3-SAT to Independent-Set

3-SAT to Independent-Set

Given a 3-CNF formula, convert into a graph

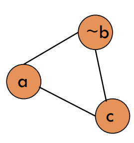
For the boolean formula in 3-SAT to be satisfied, at least one of the literals in each clause must be true

In addition, we must make sure that we enforce a literal and its complement must not both be true.

3-SAT to Independent-Set

Given a 3-CNF formula, convert into a graph

For each clause, e.g. *(a OR not(b) OR C)* create a clique containing vertices representing these literals

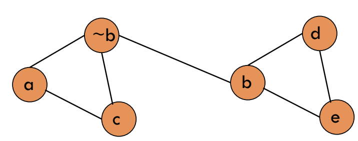


- for the Independent-Set problem to be satisfied it can only select one variable
- to make sure that all clauses are satisfied, we set k = number of clauses

3-SAT to Independent-Set

Given a 3-CNF formula, convert into a graph

To enforce that only one variable and its complement can be set we connect each vertex representing x to each vertex representing its complement ~x



Proof

Given a 3-SAT problem with k clauses and a valid truth assignment, show that f(3-SAT) has an independent set of size k. (Assume you know the solution to the 3-SAT problem and show how to get the solution to the independent set problem)

Proof

Given a 3-SAT problem with k clauses and a valid truth assignment, show that f(3-SAT) has an independent set of size k. (Assume you know the solution to the 3-SAT problem and show how to get the solution to the independent set problem)

Since each clause is an OR of variables, at least one of the three must be true for the entire formula to be true. Therefore each 3-clique in the graph will have at least on node that can be selected.

Proof

Given a graph with an independent set S of k vertices, show there exists a truth assignment satisfying the boolean formula

- For any variable xi, S cannot contain both xi and ¬xi since they are connected by an edge
- For each vertex in S, we assign it a true value and all others false. Since S has only k vertices, it must have one vertex per clause

More NP-Complete problems

SUBSET-SUM:

- Given a set S of positive integers, is there some subset S'⊆ S whose elements sum to t.

TRAVELING-SALESMAN:

- Given a weighted graph G, does the graph contain a hamiltonian cycle of length k or less?

VERTEX-COVER:

- Given a graph G = (V, E), is there a subset V'⊆V such that if (u,v)∈E then u∈V' or v∈V'?
- The extra credit was to solve this problem for bipartite graphs

Our known NP-Complete problems

We can reduce any of these problems to a new problem in an NP-completeness proof

- SAT, 3-SAT
- CLIQUE, HALF-CLIQUE
- INDEPENDENT-SET
- HAMILTONIAN-CYCLE
- TRAVELING-SALESMAN
- VERTEX-COVER
- SUBSET-SUM

Search vs. Exists

All the problems we've looked at asked decision questions:

- Is there a hamiltonian cycle?
- Does the graph have a clique of size k?
- Does the graph has an independent set of size k?
- …

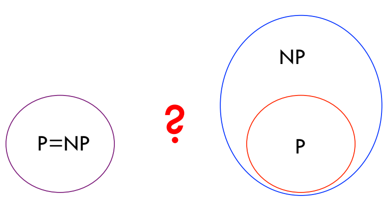
For many of the problems with a k in them, we really want to know what the largest/smallest one is

- What is the largest clique in the graph?
- What is the shortest path that visits all the vertices exactly once?

Why don't we care?

P vs. NP

The big question:



Someone finds a polynomial time solution to one of the NP-Complete problems

NP-Complete problems are somehow harder and distinct

Solving NP-Complete problems

http://www.tsp.gatech.edu/index.html